

Cryptographic Function Identification in Obfuscated Binary Programs

Hackito Ergo Sum 2012

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Presentation Outline

- Introduction to the Problem
- Proposed Solution
- Examples
- What's Next ?

INTRODUCTION TO THE PROBLEM

```

00401020 push    ebp
00401021 mov     ebp, esp
00401023 sub     esp, 14h
00401026 mov     [ebp+var_8], 20h
0040102D mov     eax, [ebp+v]
00401030 mov     ecx, [eax]
00401032 mov     [ebp+var_4], ecx
00401035 mov     edx, [ebp+v]
00401038 mov     eax, [edx+4]
0040103B mov     [ebp+var_C], eax
0040103E mov     [ebp+var_14], 9E3779B9h
00401045 mov     ecx, [ebp+var_14]
00401048 shl     ecx, 5
0040104B mov     [ebp+var_10], ecx

```

```

0040104E loc_40104E:
0040104E mov     edx, [ebp+var_8]
00401051 mov     eax, [ebp+var_8]
00401054 sub     eax, 1
00401057 mov     [ebp+var_8], eax
0040105A test    edx, edx
0040105C jbe     short loc_4010BC

```

```

0040105E mov     ecx, [ebp+var_4]
00401061 shl     ecx, 4
00401064 mov     edx, [ebp+arg_4]
00401067 add     ecx, [edx+8]
0040106A mov     eax, [ebp+var_4]
0040106D add     eax, [ebp+var_10]
00401070 xor     ecx, eax
00401072 mov     edx, [ebp+var_4]
00401075 shr     edx, 5
00401078 mov     eax, [ebp+arg_4]
0040107B add     edx, [eax+0Ch]
0040107E xor     ecx, edx
00401080 mov     edx, [ebp+var_C]
00401083 sub     edx, ecx
00401085 mov     [ebp+var_C], edx
00401088 mov     eax, [ebp+var_C]
0040108B shl     eax, 4
0040108E mov     ecx, [ebp+arg_4]
00401091 add     eax, [ecx]
00401093 mov     edx, [ebp+var_C]
00401096 add     edx, [ebp+var_10]
00401099 xor     eax, edx
0040109B mov     ecx, [ebp+var_C]
0040109E shr     ecx, 5
004010A1 mov     edx, [ebp+arg_4]
004010A4 add     ecx, [edx+4]
004010A7 xor     eax, ecx
004010A9 mov     ecx, [ebp+var_4]
004010AC sub     ecx, eax
004010AE mov     [ebp+var_4], ecx
004010B1 mov     edx, [ebp+var_10]
004010B4 sub     edx, [ebp+var_14]
004010B7 mov     [ebp+var_10], edx
004010BA jnp     short loc_40104E

```

```

004010BC loc_4010BC:
004010BC mov     eax, [ebp+v]
004010BF mov     ecx, [ebp+var_4]
004010C2 mov     [eax], ecx
004010C4 mov     edx, [ebp+v]
004010C7 mov     eax, [ebp+var_C]
004010CA mov     [edx+4], eax
004010CD mov     esp, ebp
004010CF pop     ebp
004010D0 retn
004010D0 sub_401020 endp
004010D0

```

What's this ?

```

00401020 push    ebp
00401021 mov     ebp, esp
00401023 sub     esp, 14h
00401026 mov     [ebp+var_8], 20h
0040102D mov     eax, [ebp+v]
00401030 mov     ecx, [eax]
00401032 mov     [ebp+var_4], ecx
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00401038 mov     eax, [edx+4]
0040103B mov     [ebp+var_C], eax
0040103E mov     [ebp+var_14], 9E3779B9h
00401045 mov     ecx, [ebp+var_14]
00401048 shl     ecx, 5
0040104B mov     [ebp+var_10], ecx

```

9E3779B9h

```

0040104E
0040104E loc_40104E:
0040104E mov     edx, [ebp+var_8]
00401051 mov     eax, [ebp+var_8]
00401054 sub     eax, 1
00401057 mov     [ebp+var_8], eax
0040105A test    edx, edx
0040105C jbe     short loc_4010BC

```

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0040105E mov     ecx, [ebp+var_4]
00401061 shl     ecx, 4
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00401072 mov     edx, [ebp+var_4]
00401075 shr     edx, 5
00401078 mov     eax, [ebp+arg_4]
0040107B add     edx, [eax+0Ch]
0040107E xor     ecx, edx
00401080 mov     edx, [ebp+var_C]
00401083 sub     edx, ecx
00401085 mov     [ebp+var_C], edx
00401088 mov     eax, [ebp+var_C]
0040108B shl     eax, 4
0040108E mov     ecx, [ebp+arg_4]
00401091 add     eax, [ecx]
00401093 mov     edx, [ebp+var_C]
00401096 add     edx, [ebp+var_10]
00401099 xor     eax, edx
0040109B mov     ecx, [ebp+var_C]
0040109E shr     ecx, 5
004010A1 mov     edx, [ebp+arg_4]
004010A4 add     ecx, [edx+4]
004010A7 xor     eax, ecx
004010A9 mov     ecx, [ebp+var_4]
004010AC sub     ecx, eax
004010AE mov     [ebp+var_4], ecx
004010B1 mov     edx, [ebp+var_10]
004010B4 sub     edx, [ebp+var_14]
004010B7 mov     [ebp+var_10], edx
004010BA jnp     short loc_40104E

```

```

004010BC
004010BC loc_4010BC:
004010BC mov     eax, [ebp+v]
004010BF mov     ecx, [ebp+var_4]
004010C2 mov     [eax], ecx
004010C4 mov     edx, [ebp+v]
004010C7 mov     eax, [ebp+var_C]
004010CA mov     [edx+4], eax
004010CD mov     esp, ebp
004010CF pop     ebp
004010D0 retn
004010D0 sub_401020 endp
004010D0

```

What's this ?

[Tiny Encryption Algorithm - Wikipedia, the free encyclopedia](#)

en.wikipedia.org/wiki/Tiny_Encryption_Algorithm - Traduire cette page

The magic constant, 2654435769 or **9E3779B9**16 is chosen to be $232/\phi$, where ϕ is the golden ratio. TEA has a few weaknesses. Most notably, it suffers from ...

↳ [Properties](#) - [Versions](#) - [Reference code](#) - [See also](#)

Vous avez consulté cette page 15 fois. Dernière visite : 07/04/12

[PDF\] The RC5 Encryption Algorithm?](#)

www.engr.uconn.edu/~zshi/.../rc5.pdf - États-Unis - Traduire cette page

Format de fichier: PDF/Adobe Acrobat - [Afficher](#)

de RL Rivest - [Cité 827 fois](#) - [Autres articles](#)

Q32 = 10011110001101110111100110111001 = **9e3779b9**. P64 =
1011011111100001010100010110001010001010111011010010101001101011 ...

Vous avez consulté cette page le 05/04/12.

[Internet Security: Cryptographic Principles, Algorithms and Protocols - Résultats Google Recherche de Livres](#)

books.google.fr/books?isbn=0470852852...

Man Young Rhee - 2003 - Computers - 405 pages

... b7e!5163 + **9e3779b9** = 5618cblc 5[2] = 5[1] + Q32 = 5618cblc + **9e3779b9** =
f45044d5 5[3] = S[2] + Q32 = f45044d5 + **9e3779b9** = 9287be8e S[25] = S[24] + ...

[Changeset 329 – CrypTool 2.0](#)

<https://www.cryptool.org/trac/CrypTool2/changeset/329>

28 May 2009 – The magic constant, 2654435769 (Decimal) or **9E3779B9** (Hex) is chosen to be $(2^{32} / \phi)$ where ϕ is the golden ratio. </Run> ...

[TEA Encryption Algorithm, Source code](#)

www.shokhirev.com/nikolai/.../uTeaSet_pas.html - Traduire cette page

... XTeaEncrypt/XTeaDecrypt Thanks to Pedro Gimeno Fortea <parigalo@formauri.es>

```
} interface const Delta: longword = $9e3779b9; type TLong2 = array[0.
```

Tools	Answer
Crypto Searcher	"TEA"
Draca v0.5.7b	"TEA/RC5/RC6"
Findcrypt v2	∅
Hash & Crypto Detector v1.4	"TEA/XTEA/TEAN"
PEiD KANAL v2.92	"TEA/N, RC5, RC6"
Kerckhoffs	∅
Signsrch 0.1.7	"TEA"
SnD Crypto Scanner v0.5b	∅

Tools	Answer
Crypto Searcher	“TEA”
Draca v0.5.7b	“TEA/RC5/RC6”
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Hash & Crypto Detector v1.4	“TEA/XTEA/TEAN”
PEiD KANAL v2.92	“TEA/N, RC5, RC6”
Kerckhoffs	∅
Signsrch 0.1.7	“TEA”
SnD Crypto Scanner v0.5b	∅

That’s indeed the Tiny Encryption Algorithm!


```

nov    [ebp+var_10], ecx
nov    edx, [ebp+arg_C]
imul  edx, [ebp+arg_8]
nov    [ebp+var_14], edx
nov    eax, [ebp+arg_4]
nov    ecx, [eax]
nov    [ebp+var_8], ecx
nov    edx, [ebp+arg_4]
nov    eax, [edx+4]
nov    [ebp+var_20], eax
nov    ecx, [ebp+arg_4]
nov    edx, [ecx+8]
nov    [ebp+var_1C], edx
nov    eax, [ebp+arg_4]
nov    ecx, [eax+0Ch]
nov    [ebp+var_18], ecx
nov    [ebp+var_4], 0
jmp    short loc_40105F

```

```

loc_40105F:
cmp    [ebp+var_4], 20h
jnb    short loc_401088

```

```

nov    eax, [ebp+var_24]
shl   eax, 4
add   eax, [ebp+var_1C]
nov    ecx, [ebp+var_24]
add   ecx, [ebp+var_14]
xor   eax, ecx
nov    edx, [ebp+var_24]
shr   edx, 5
add   edx, [ebp+var_18]
xor   eax, edx
nov    ecx, [ebp+var_C]
sub   ecx, eax
nov    [ebp+var_C], ecx
nov    edx, [ebp+var_C]
shl   edx, 4
add   edx, [ebp+var_8]
nov    eax, [ebp+var_C]
add   eax, [ebp+var_14]
xor   edx, eax
nov    ecx, [ebp+var_C]
shr   ecx, 5
add   ecx, [ebp+var_20]
xor   edx, ecx
nov    eax, [ebp+var_24]
sub   eax, edx
nov    [ebp+var_24], eax
nov    ecx, [ebp+var_14]
sub   ecx, [ebp+var_10]
nov    [ebp+var_14], ecx
jmp    short loc_401056

```

```

loc_401088:
nov    edx, [ebp+arg_0]
nov    eax, [ebp+var_24]
mov    [edx], eax
nov    ecx, [ebp+arg_0]
nov    edx, [ebp+var_C]
mov    [ecx+4], edx
nov    esp, ebp
pop    ebp
retn
?decrypt@YAXPAI0HIGZ endp

```

```

loc_401056:
nov    edx, [ebp+var_4]
add   edx, 1
nov    [ebp+var_4], edx

```

What about this one?

```

mov [ebp+var_10], ecx
mov edx, [ebp+arg_C]
imul edx, [ebp+arg_8]
mov [ebp+var_14], edx
mov eax, [ebp+arg_4]
mov ecx, [eax]
mov [ebp+var_8], ecx
mov edx, [ebp+arg_4]
mov eax, [edx+4]
mov [ebp+var_20], eax
mov ecx, [ebp+arg_4]
mov edx, [ecx+8]
mov [ebp+var_1C], edx
mov eax, [ebp+arg_4]
mov ecx, [eax+0Ch]
mov [ebp+var_18], ecx
mov [ebp+var_4], 0
jmp short loc_40105F

```



```

mov [ebp+var_10], ecx
mov edx, [ebp+arg_C]
imul edx, [ebp+arg_8]
mov [ebp+var_14], edx
mov eax, [ebp+arg_4]
mov ecx, [eax]
mov [ebp+var_8], ecx
mov edx, [ebp+arg_4]
mov eax, [edx+4]
mov [ebp+var_20], eax
mov ecx, [ebp+arg_4]
mov edx, [ecx+8]
mov [ebp+var_1C], edx
mov eax, [ebp+arg_4]
mov ecx, [eax+0Ch]
mov [ebp+var_18], ecx
mov [ebp+var_4], 0
jmp short loc_40105F

```

```

loc_40105F:
cmp [ebp+var_4], 20h
jnb short loc_401088

```

```

mov eax, [ebp+var_24]
shl eax, 4
add eax, [ebp+var_1C]
mov ecx, [ebp+var_24]
add ecx, [ebp+var_14]
xor eax, ecx
mov edx, [ebp+var_24]
shr edx, 5
add edx, [ebp+var_18]
xor eax, edx
mov ecx, [ebp+var_C]
sub ecx, eax
mov [ebp+var_C], ecx
mov edx, [ebp+var_C]
shl edx, 4
add edx, [ebp+var_8]
mov eax, [ebp+var_C]
add eax, [ebp+var_14]
xor edx, eax
mov ecx, [ebp+var_C]
shr ecx, 5
add ecx, [ebp+var_20]
xor edx, ecx
mov eax, [ebp+var_24]
sub eax, edx
mov [ebp+var_24], eax
mov ecx, [ebp+var_14]
sub ecx, [ebp+var_10]
mov [ebp+var_14], ecx
jmp short loc_401056

```

```

loc_401088:
mov edx, [ebp+arg_0]
mov eax, [ebp+var_24]
mov [edx], eax
mov ecx, [ebp+arg_0]
mov edx, [ebp+var_C]
mov [ecx+4], edx
mov esp, ebp
pop ebp
retn
?decrypt@YAXPAI0HIGZ endp

```

```

loc_401056:
mov edx, [ebp+var_4]
add edx, 1
mov [ebp+var_4], edx

```

What about this one?

No particular constants

Tools	Answer
Crypto Searcher	∅
Draca v0.5.7b	∅
Findcrypt v2	∅
Hash & Crypto Detector v1.4	∅
PEiD KANAL v2.92	∅
Kerckhoffs	∅
Signsrch 0.1.7	∅
SnD Crypto Scanner v0.5b	∅

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PEiD KANAL v2.92	∅
Kerckhoffs	∅
Signsrch 0.1.7	∅
SnD Crypto Scanner v0.5b	∅

Sigh.. That was still TEA!

What Can We Do ?

- How to recognize different TEA implementations in a more reliable way ?
- Is there something such implementations *have to* share ?

Input-Output Relationship

- For a key \mathbf{K} and an encrypted text \mathbf{C} , *any* TEA implementation produces the *same* decrypted text \mathbf{C}' .

Input-Output Relationship

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Could we identify TEA implementations by using their I/O relationship ?

Input-Output Relationship

- For a key \mathbf{K} and an encrypted text \mathbf{C} , *any* TEA implementation produces the *same* decrypted text \mathbf{C}' .

Could we identify TEA implementations by using their I/O relationship ?

(or any other cipher)

PROPOSED SOLUTION

How To Use Input-Output Relationship ?

- Let's say **P** is a program implementing an unknown cryptographic algorithm.

How To Use Input-Output Relationship ?

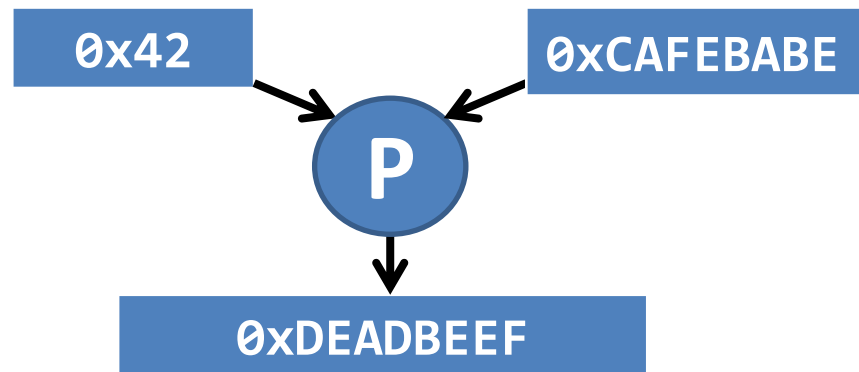
- Let's say **P** is a program implementing an unknown cryptographic algorithm.
- We can not *realistically* use the I/O relationship to prove that **P** implements a particular crypto algorithm *on any inputs*.

(too many input states to test!)

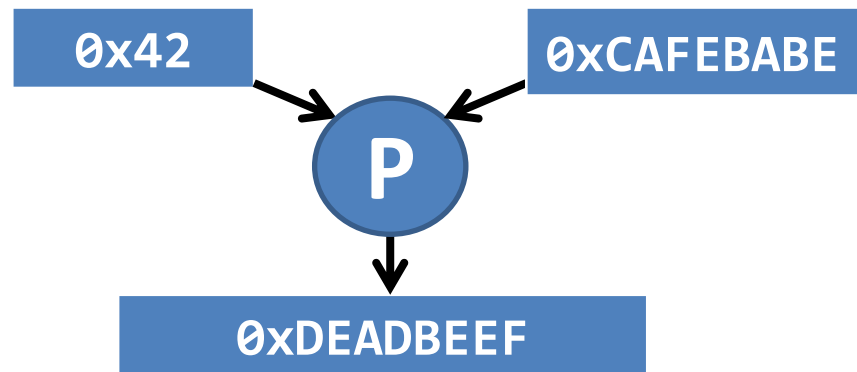
How To Use Input-Output Relationship ?

- Let's say **P** is a program implementing an unknown cryptographic algorithm.
- We can not *realistically* use the I/O relationship to prove that **P** implements a particular crypto algorithm *on any inputs*.
(too many input states to test!)
- But we can observe one particular **P** execution and collect its input-output values...

For example:



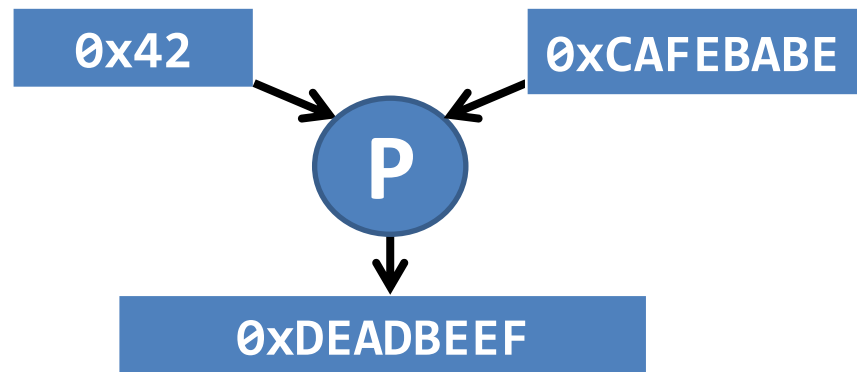
For example:



- Now imagine that when we execute a reference implementation of **TEA** with the key 0x42 and the input text 0xCAFEBAE, it produces 0xDEADBEEF.

What does it mean ?

For example:



- Now imagine that when we execute a reference implementation of **TEA** with the key **0x42** and the input text **0xCAFEBAE**, it produces **0xDEADBEEF**.

What does it mean ?

- It proves that **P** implements **TEA** on these particular input values.

Final Goal

- We are going to prove that a particular program **P** behaves like a known cryptographic algorithm during a particular execution.
- It means that we are **not** going to prove a general semantic equivalence between **P** and a cryptographic algorithm.

Workflow

Given a program **P**:

- **Step 1:** Collect **P** execution trace.
- **Step 2:** Extract possible cryptographic algorithms *with their parameters* from **P** execution trace (here is the magic).
- **Step 3:** Identify these algorithms by comparing their I/O relationship with those of known algorithms.

STEP 1: COLLECT EXECUTION TRACE

Execution Trace

- Pin: Dynamic Binary Instrumentation framework.

Address	Instruction	Read Registers	Written Registers	Read Memory	Written Memory
4012b3	push ebp	ebp 0012de28 esp 0012bd98	esp 0012bd94		12bd94 0012de28
4012b4	mov ebp, esp	esp 0012bd94	ebp 0012bd94		
4012b6	push ebx	ebx 02f00010 esp 0012bd84	esp 0012bd80		12bd80 2f00010

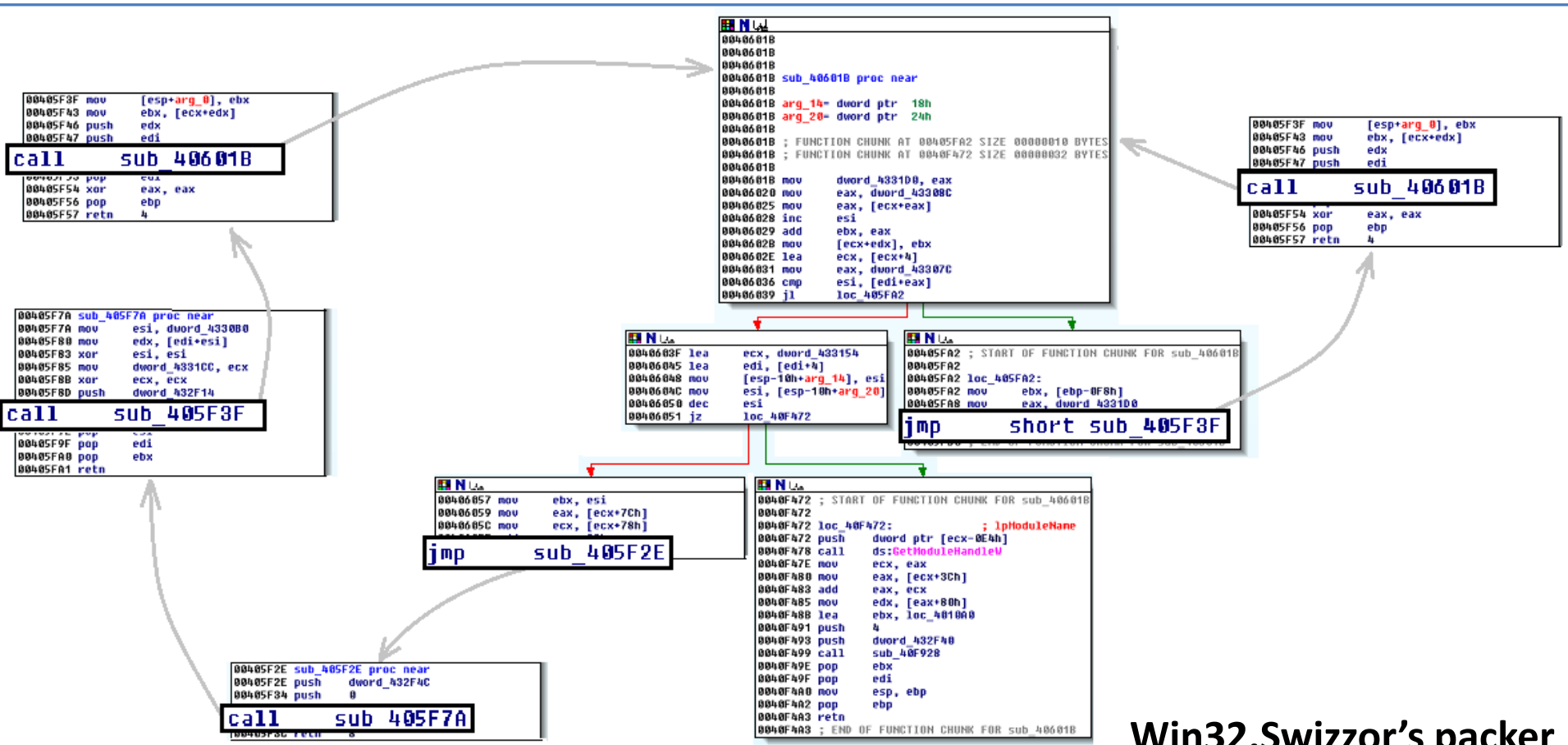
...

STEP 2: CRYPTOGRAPHIC ALGORITHM EXTRACTION

How To Find Crypto Code ? (1)

- Cryptographic code constitutes only a part of programs, we need a way to find it.
- As we want to play with *obfuscated* programs, IDA functions will not be enough...

In obfuscated programs, such things can happen:



Win32.Swizzor's packer

How To Find Crypto Code ? (2)

How To Find Crypto Code ? (2)

- Cryptographic algorithms usually apply *a same treatment* on their input-output parameters.

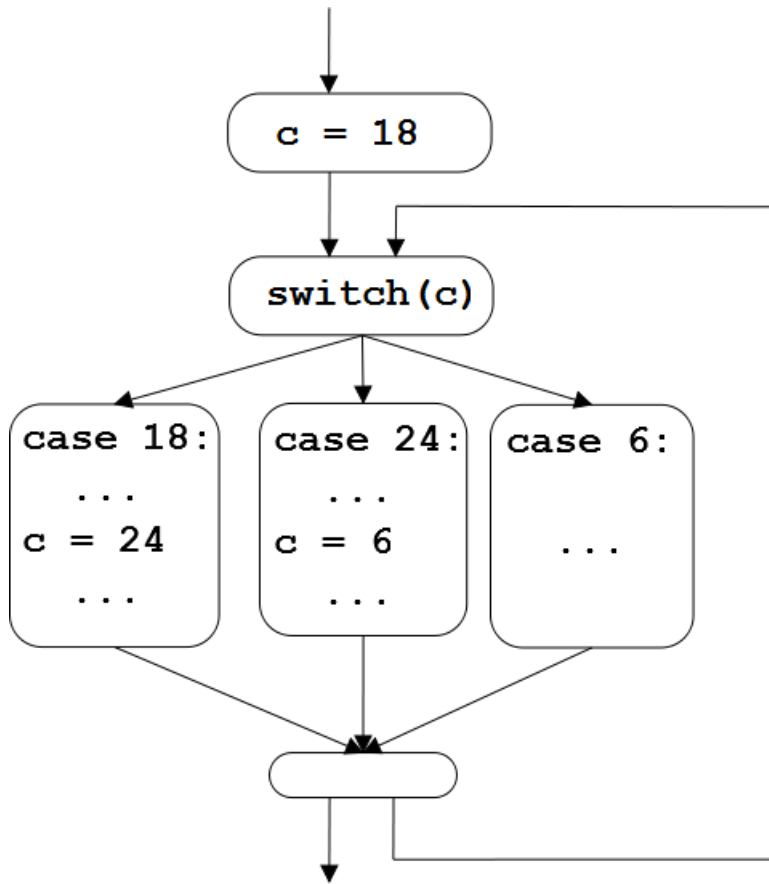
How To Find Crypto Code ? (2)

- Cryptographic algorithms usually apply *a same treatment* on their input-output parameters.
- It makes **loops** a cryptographic code feature.

How To Find Crypto Code ? (2)

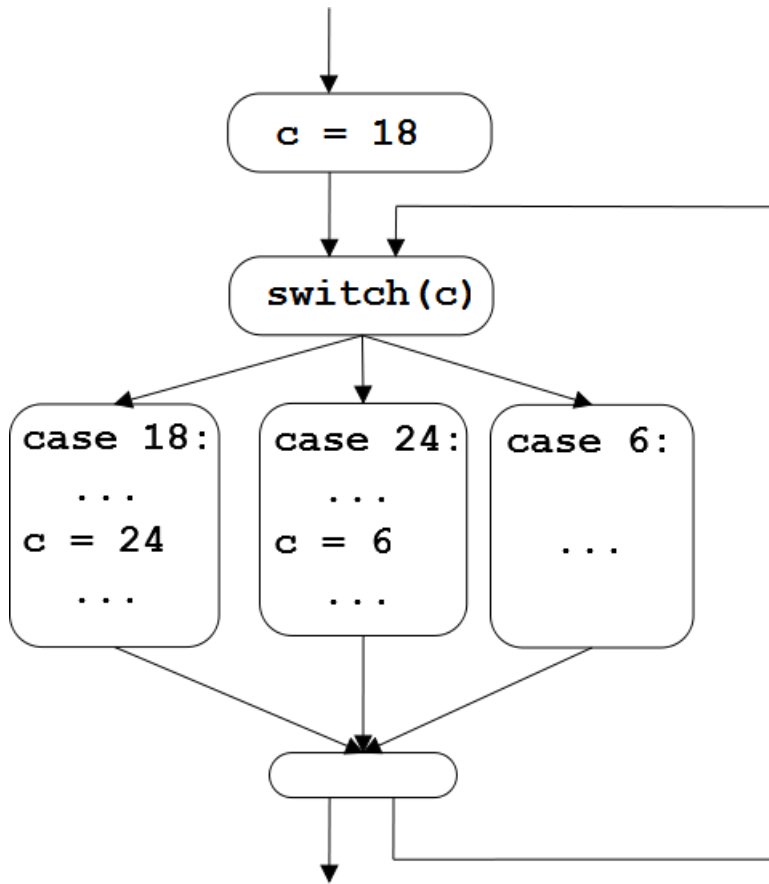
- Cryptographic algorithms usually apply *a same treatment* on their input-output parameters.
- It makes **loops** a cryptographic code feature.
- But there are loops everywhere, not only in crypto... What kind of loops are we looking for ?

Loops ?

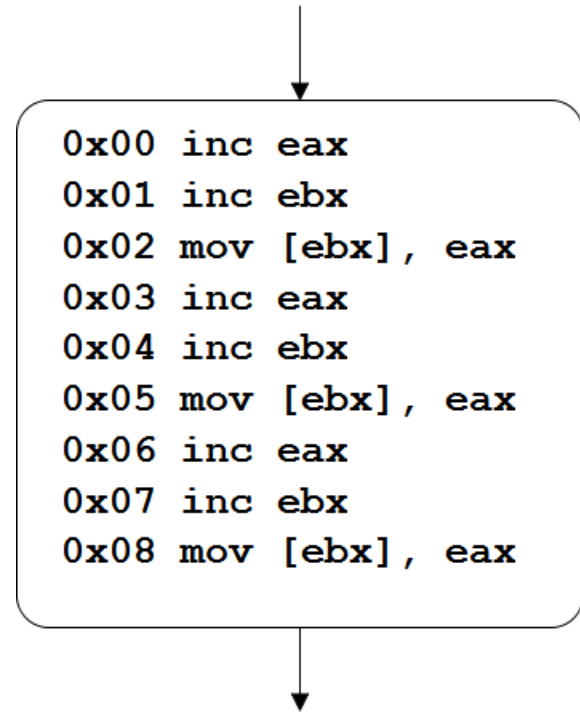


Mebroot state-machine

Loops ?

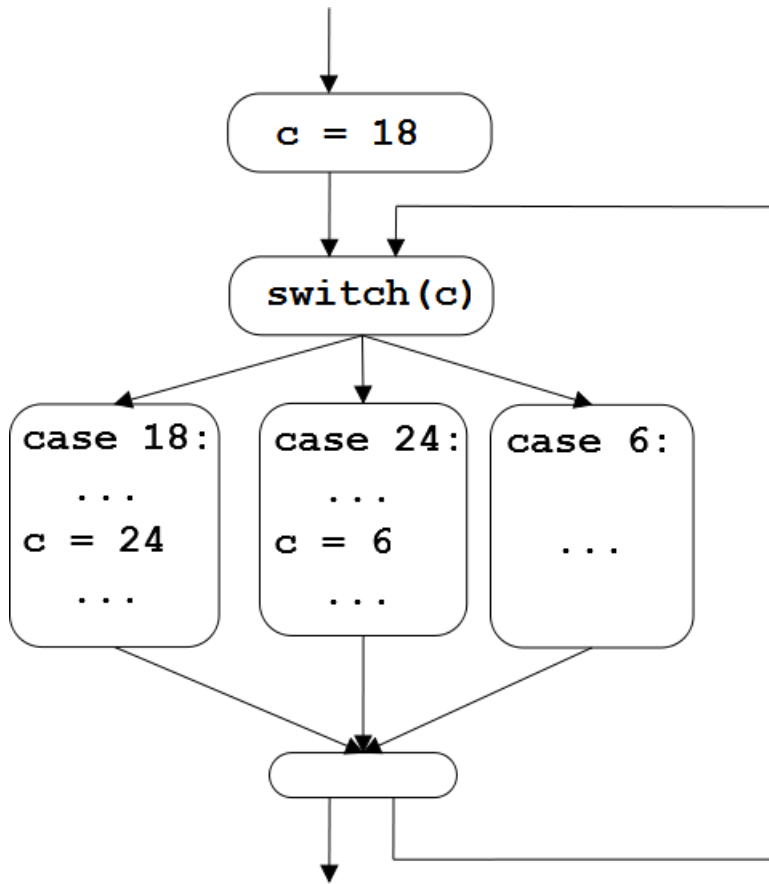


Mebroot state-machine

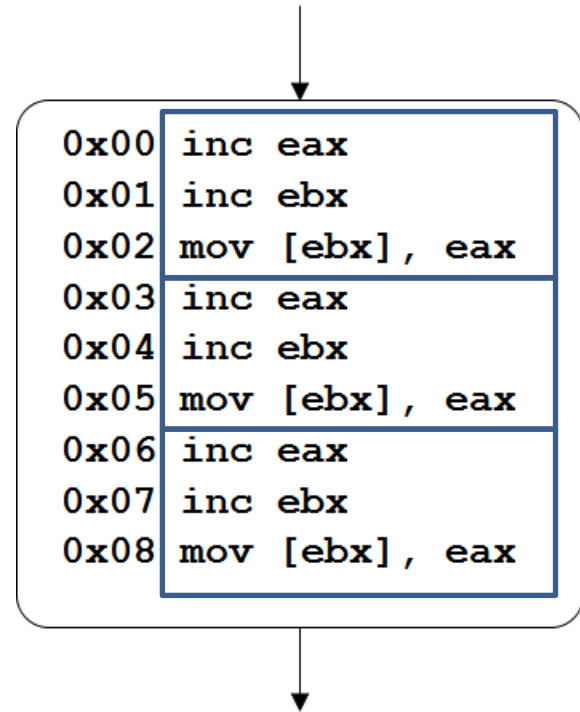


Unrolling optimization

Loops ?



Mebroot state-machine



Unrolling optimization

Loooooops

- We look for **the same operations applied repeatedly** on a set of data.

Loooooops

- We look for **the same operations applied repeatedly** on a set of data.

“A loop is the repetition of a same sequence of machine instructions at least two times.”

(This sequence of instructions is the loop body.)

Example

Execution Trace

...	...
401325	add ebx, edi
401327	sub edx, ebx
401329	dec dword ptr [ebp+0xc]
40132c	jnz 0x401325
401325	add ebx, edi
401327	sub edx, ebx
401329	dec dword ptr [ebp+0xc]
40132c	jnz 0x401325
...	...

Example

Execution Trace

...	...	
401325	add ebx, edi	Iteration 1
401327	sub edx, ebx	
401329	dec dword ptr [ebp+0xc]	
40132c	jnz 0x401325	
401325	add ebx, edi	Iteration 2
401327	sub edx, ebx	
401329	dec dword ptr [ebp+0xc]	
40132c	jnz 0x401325	
...	...	

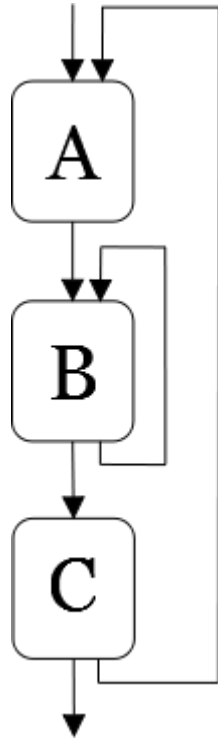
Example

Execution Trace

...	...	
401325	add ebx, edi	Iteration 1
401327	sub edx, ebx	
401329	dec dword ptr [ebp+0xc]	
40132c	jnz 0x401325	
401325	add ebx, edi	Iteration 2
401327	sub edx, ebx	
401329	dec dword ptr [ebp+0xc]	
40132c	jnz 0x401325	
...	...	

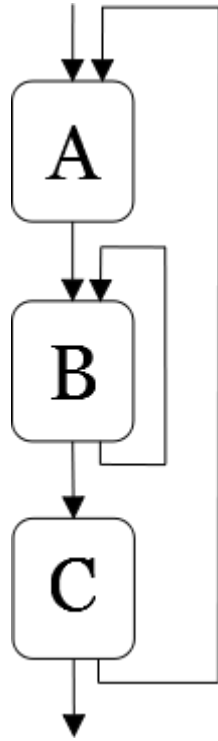
Loop

What About Nested Loops ?



Simplified CFG

What About Nested Loops ?

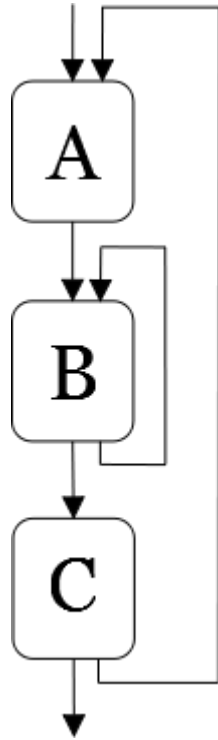


Simplified CFG

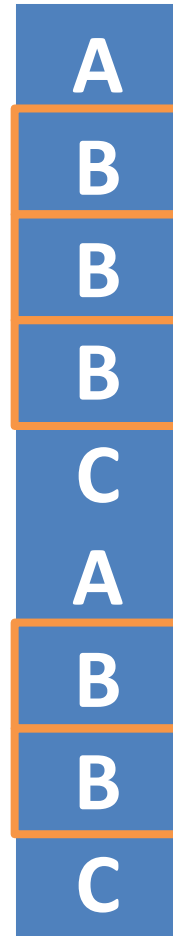


Execution trace

What About Nested Loops ?



Simplified CFG

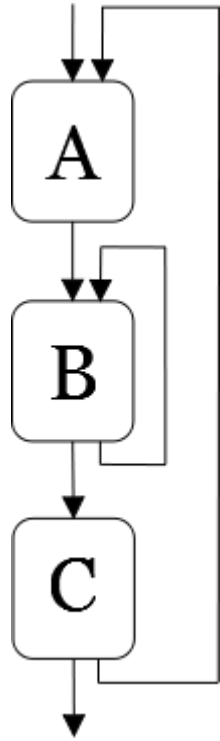


Loop B
3 iterations

Loop B
2 iterations

Execution trace

What About Nested Loops ?

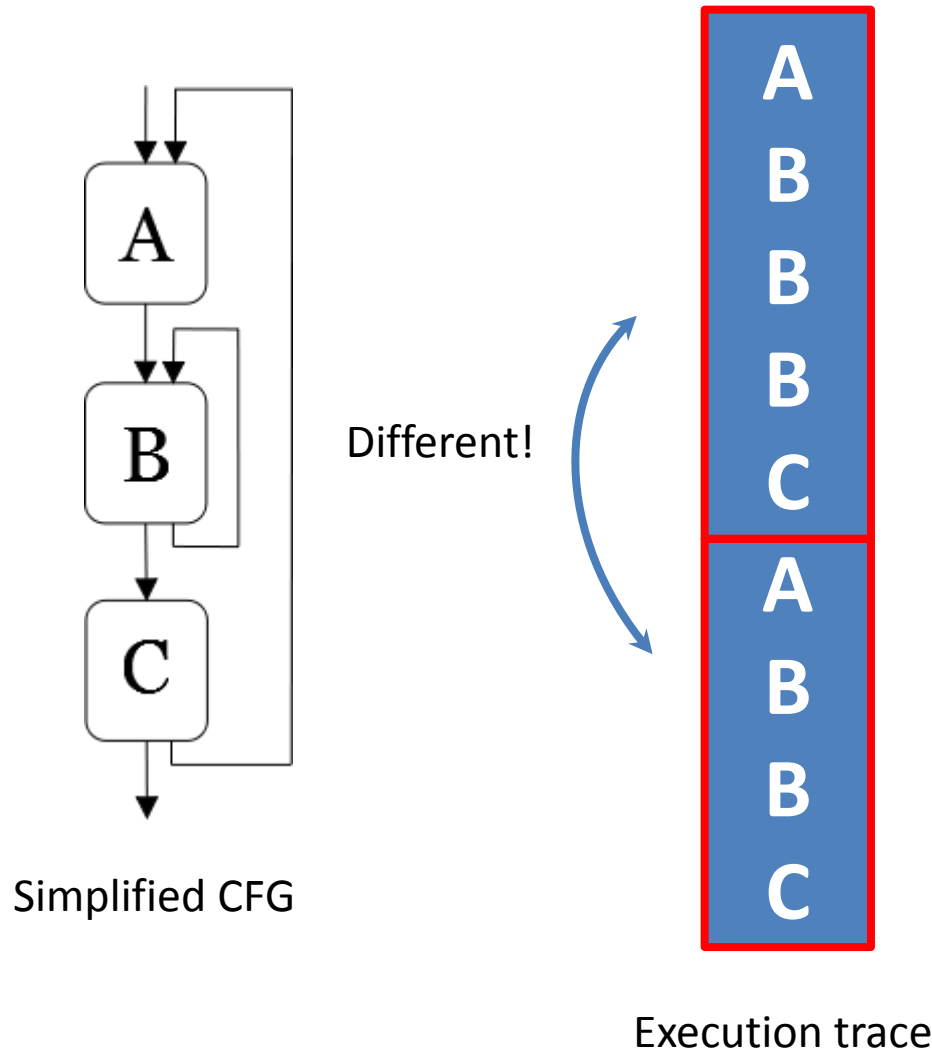


Simplified CFG

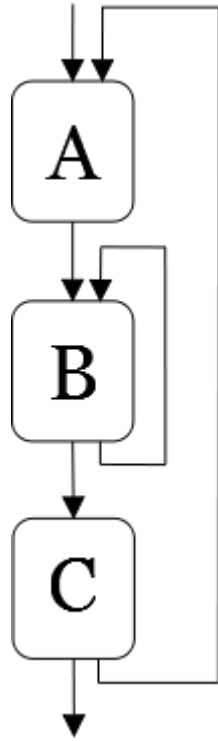


Execution trace

What About Nested Loops ?



What About Nested Loops ?

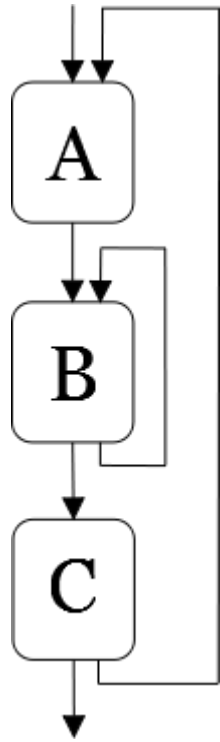


Simplified CFG

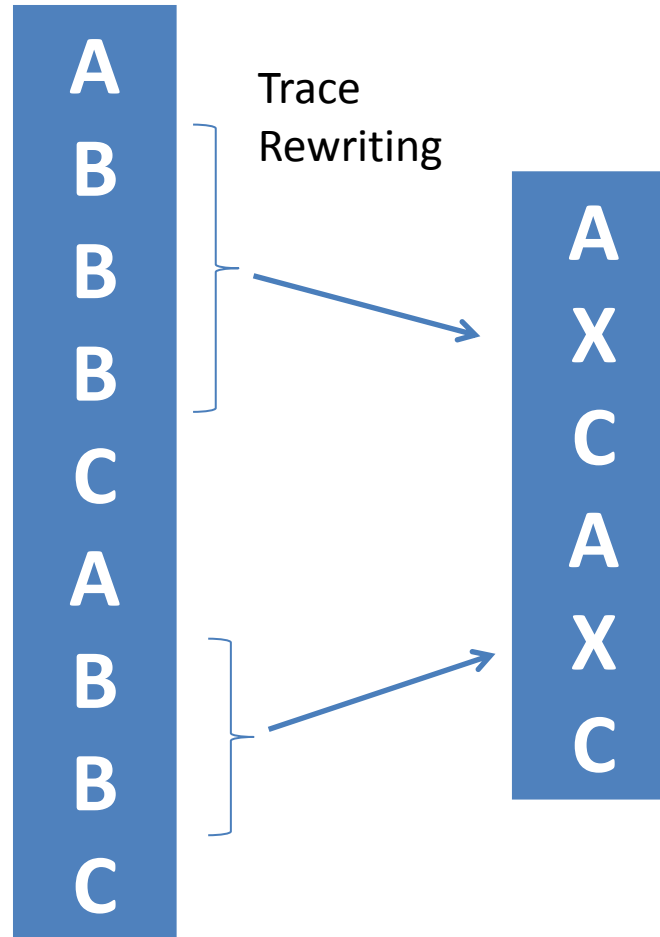


Execution trace

What About Nested Loops ?



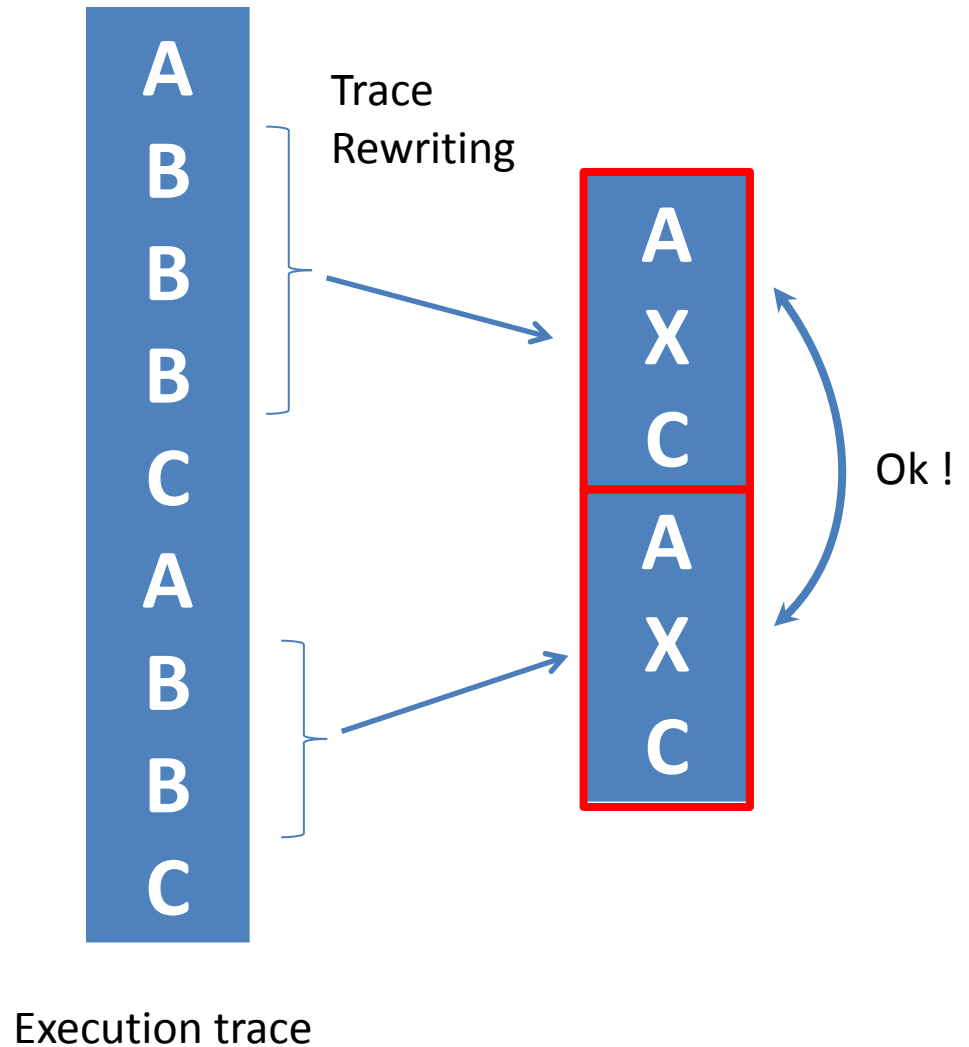
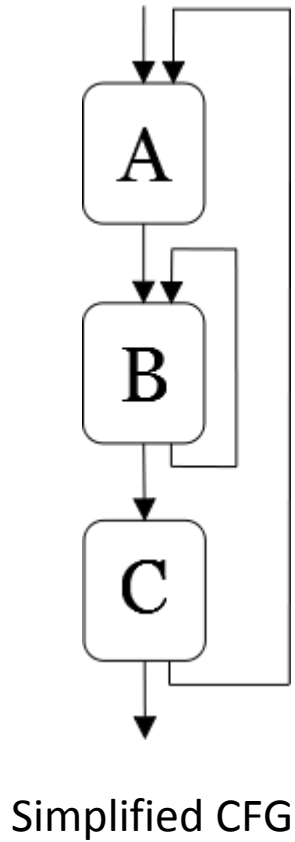
Simplified CFG



Execution trace

Ok !

What About Nested Loops ?



Loop Detection Algorithm

1. Detects two repetitions of a loop body in the execution trace.

(non trivial, because there is an infinity of possible loop body)

2. Replaces in the trace the detected loop by a symbol representing their body.
3. Goes back to step 1 if new loops have been detected.

What's Next ?

- We extracted possible cryptographic code from execution traces thanks to a particular loop definition.

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- We extracted possible cryptographic code from execution traces thanks to a particular loop definition.
- For the moment, we assume that each possible cryptographic algorithm corresponds to *one single* loop.
- How can we define *parameters* from the *bytes* read and written in the execution trace ?

Loop Parameters (1)

- Distinction between input and output bytes in the execution trace:
 - **Input** bytes have been **read without having been previously written.**
 - **Output** bytes have been **written.**

Loop Parameters (2)

- We want to group together bytes belonging to the same cryptographic parameter (key, input text...).

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What criteria can we use ?

Loop Parameters (3)

- Grouping of several bytes into the same parameter:
 1. If they are **adjacent in memory** (*too large!*)

Loop Parameters (3)

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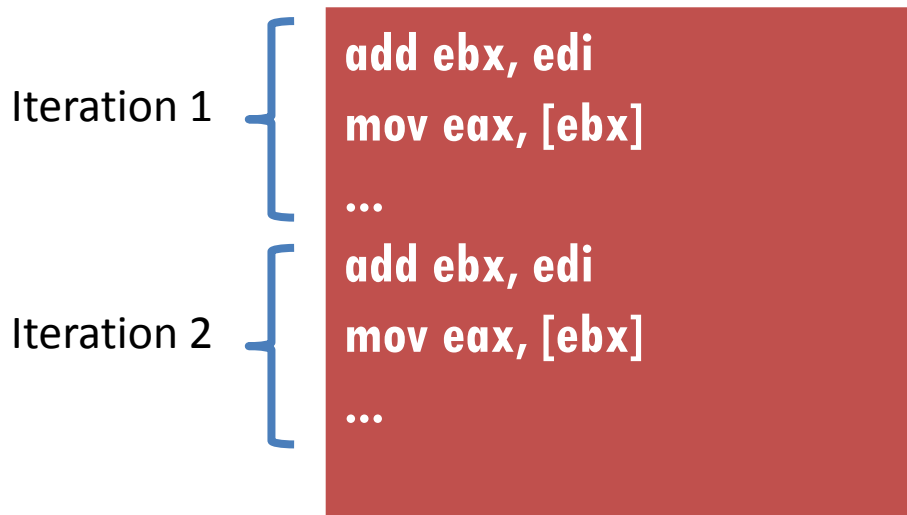
Loop Parameters (3)

- Grouping of several bytes into the same parameter:
 1. If they are **adjacent in memory** (*too large!*)
 2. And if they are **manipulated by a same instruction in the loop body.**

```
add ebx, edi
mov eax, [ebx]
...
add ebx, edi
mov eax, [ebx]
...
```

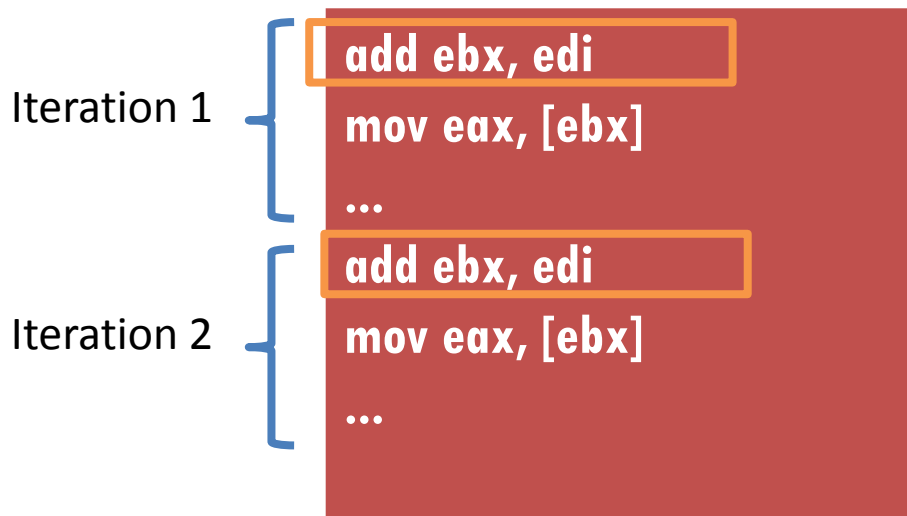
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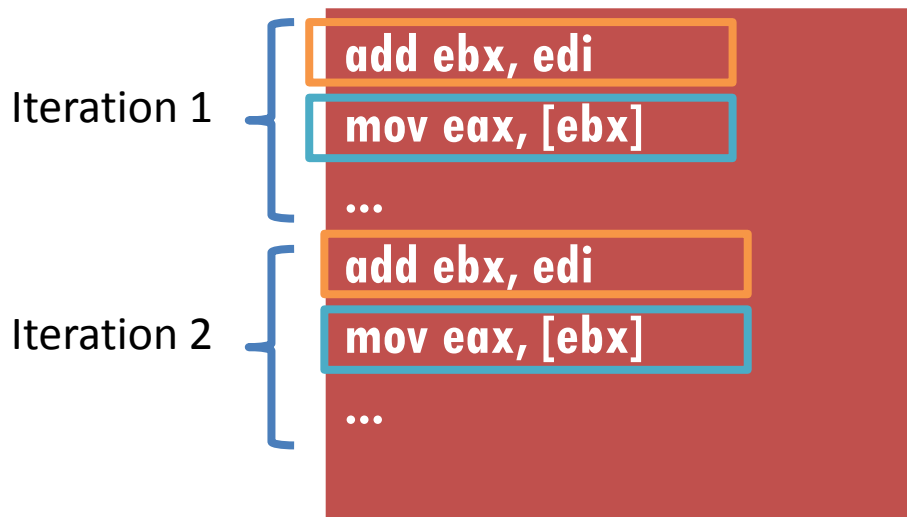
Loop Parameters (3)

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Loop Parameters (3)

- Grouping of several bytes into the same parameter:
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Loop Parameters (4)

- Once a parameter has been defined, we collect its value from the trace.

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- Once a parameter has been defined, we collect its value from the trace.

Parameter

`(memory address|register name):size`

Value in hexadecimal

Let's Recap With a Use-Case

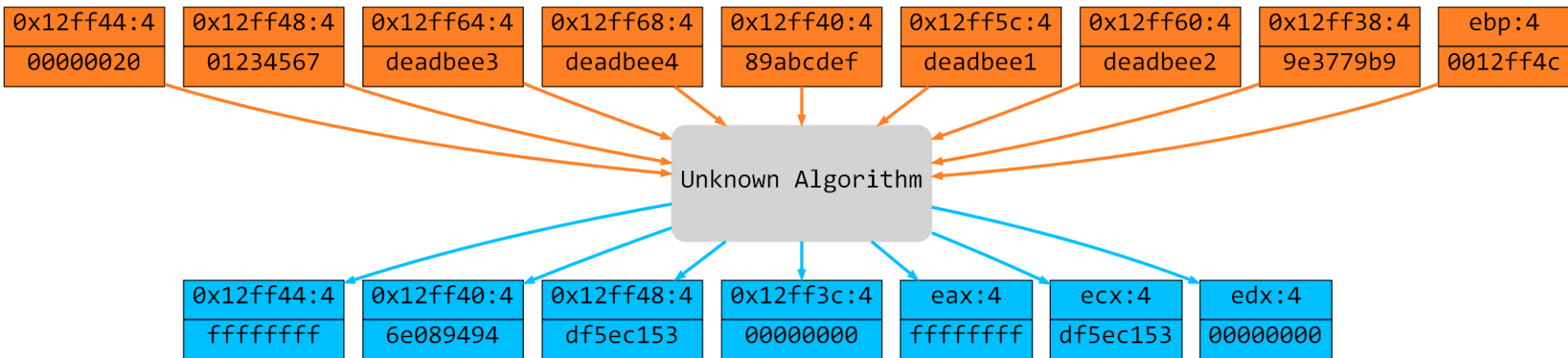
- Tiny Encryption Algorithm:
 - Block cipher
 - 64 round Feistel network
 - 16-byte key
 - 8-byte input text
 - Magic constant ***delta*** (0x9E3779B9)
- We built a toy program calling the TEA decryption function on:
 - Key : **0xDEADBEE1...DEADBEE4**
 - Encrypted text: **0x0123456789ABCDEF**

```
4010d0!push ebp!RR_ebp_12ffc0_esp_12ff80!WM_12ff7c_4_12ffc0!WR_esp_12ff7c
4010d1!mov ebp, esp!RR_esp_12ff7c!WR_ebp_12ff7c
4010d3!sub esp, 0x18!RR_esp_12ff7c!WR_esp_12ff64
4010d6!mov dword ptr [ebp-0x8], 0x01234567!RR_ebp_12ff7c!WM_12ff74_4_cafebabe
4010dd!mov dword ptr [ebp-0x4], 0x890ABCDEF!RR_ebp_12ff7c!WM_12ff78_4_cafebabe
4010e4!mov dword ptr [ebp-0x18], 0xdeadbee1!RR_ebp_12ff7c!WM_12ff64_4_deadbeef
4010eb!mov dword ptr [ebp-0x14], 0xdeadbee2!RR_ebp_12ff7c!WM_12ff68_4_deadbeef
4010f2!mov dword ptr [ebp-0x10], 0xdeadbee3!RR_ebp_12ff7c!WM_12ff6c_4_deadbeef
4010f9!mov dword ptr [ebp-0xc], 0xdeadbee4!RR_ebp_12ff7c!WM_12ff70_4_deadbeef
401100!lea eax, ptr [ebp-0x18]!RR_ebp_12ff7c!WR_eax_12ff64
```

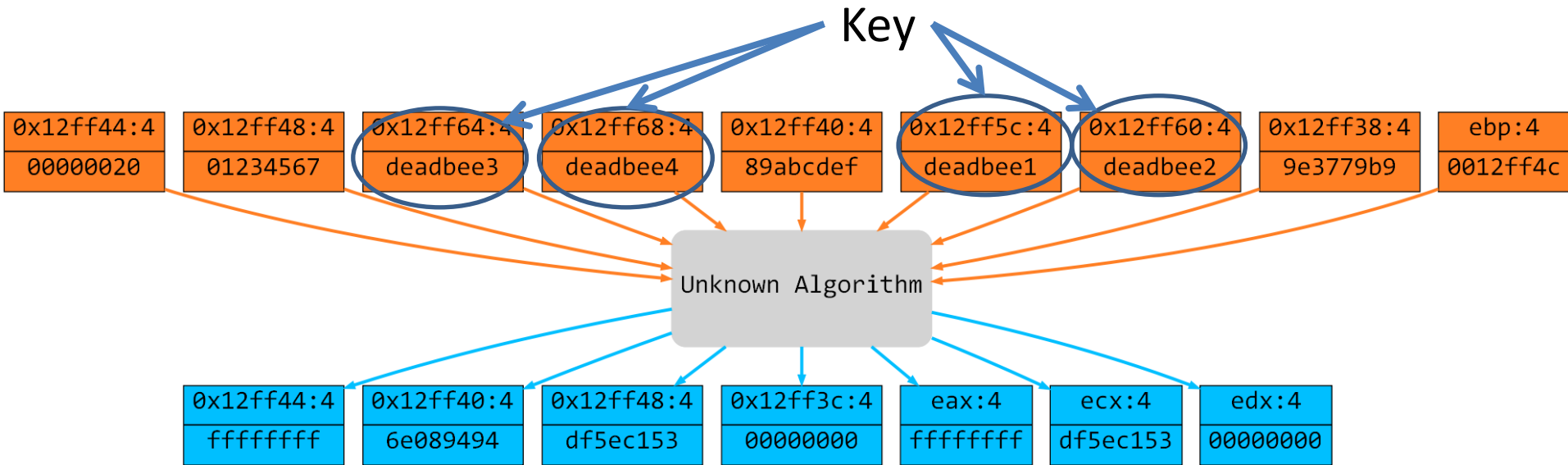
...

Step 1: collect the execution trace

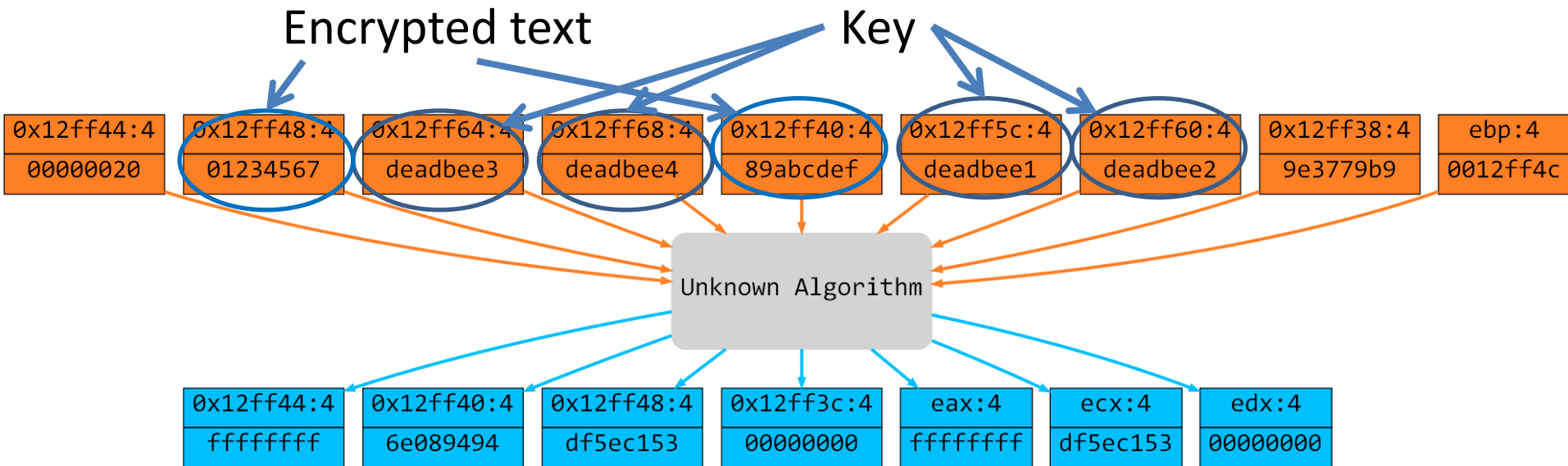
Step 2: cryptographic algorithm
extraction



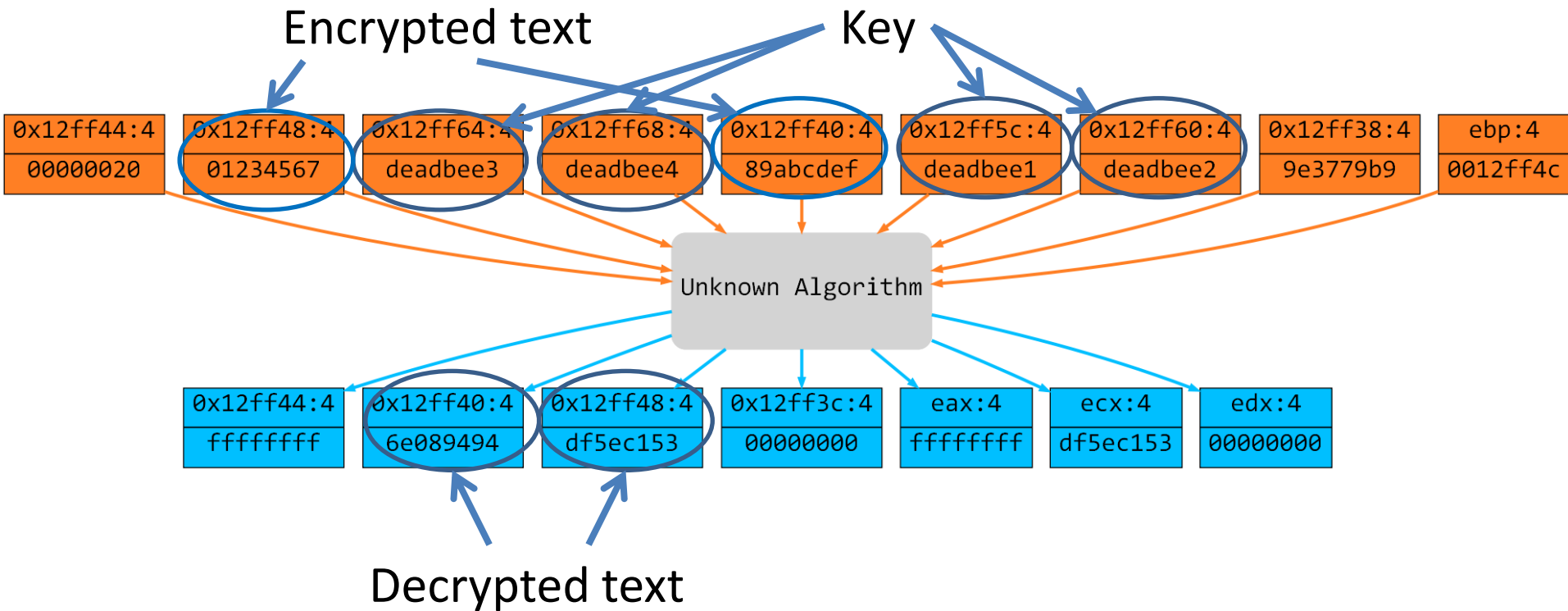
Step 2: cryptographic algorithm
extraction



Step 2: cryptographic algorithm extraction



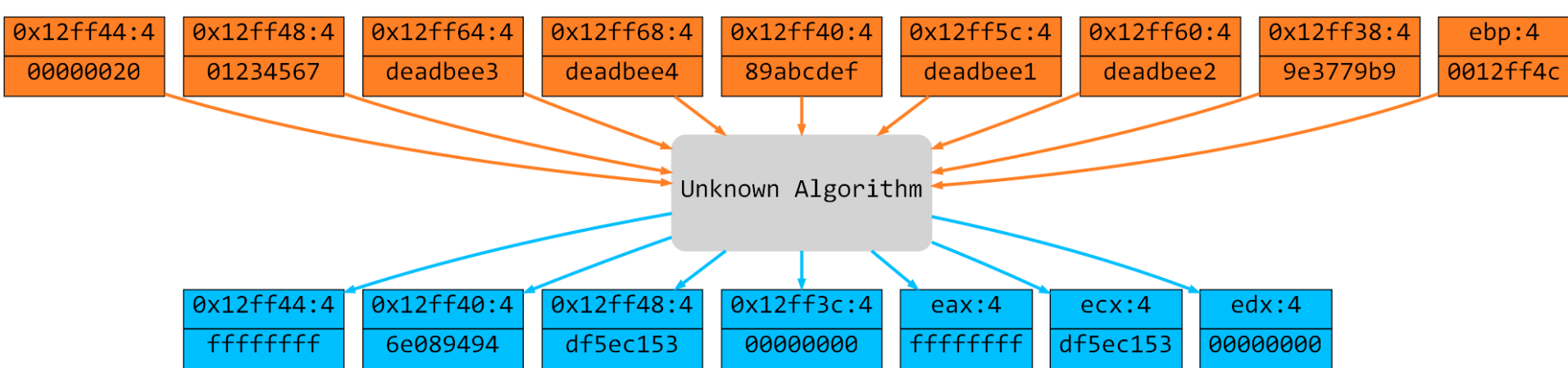
Step 2: cryptographic algorithm extraction



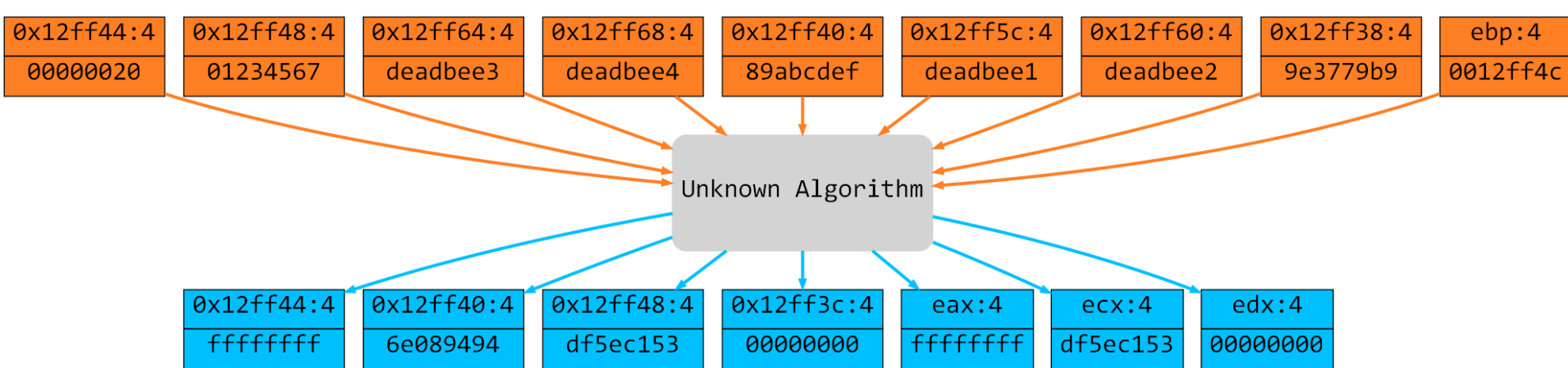
Step 2: cryptographic algorithm extraction

STEP 3: CRYPTO ALGORITHM IDENTIFICATION

Input 1: unknown algorithm **A** with its parameter values



Input 1: unknown algorithm **A** with its parameter values



Input 2: reference implementations for common crypto algo

```
def tea (input_text, key):
```

...

```
def xtea (input_text, key):
```

...

```
def rc4 (input_text, key):
```

...

Question

- Is there a way to combine **A** input values such that *tea()*, *xtea()* or *rc4()* would produce **A** output values ?

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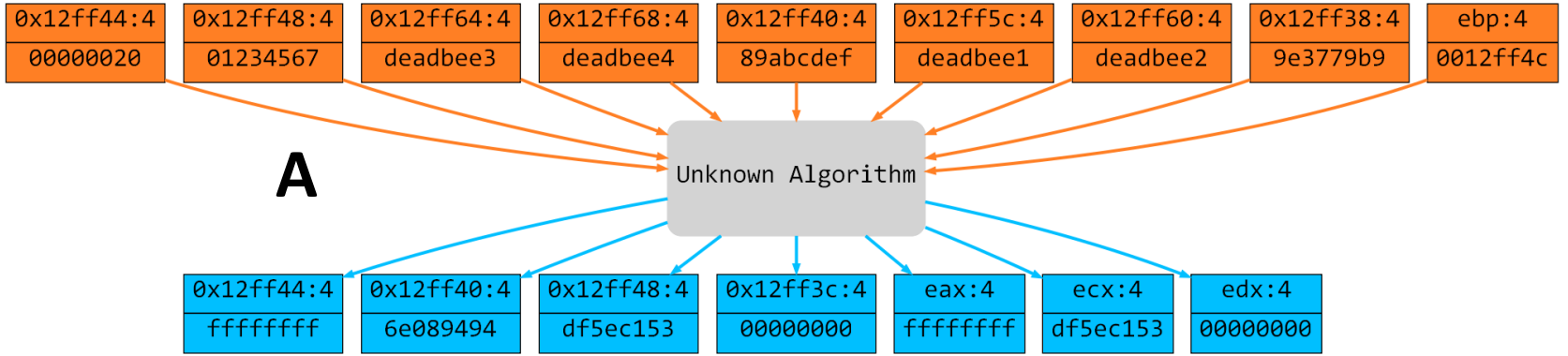
Question

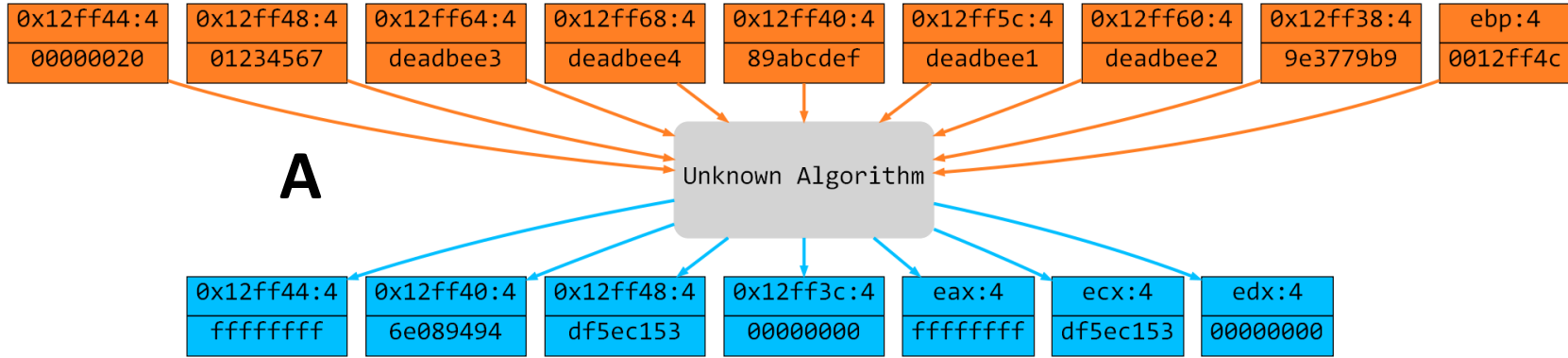
- Is there a way to combine **A** input values such that *tea()*, *xtea()* or *rc4()* would produce **A** output values ?
- Some difficulties:
 - **Parameter division**: a same cryptographic parameter can be divided into several loop parameter.
 - **Parameter order**: no particular order for **A** parameters.

Question

- Is there a way to combine **A** input values such that *tea()*, *xtea()* or *rc4()* would produce **A** output values ?
- Some difficulties:
 - **Parameter division**: a same cryptographic parameter can be divided into several loop parameter.
 - **Parameter order**: no particular order for **A** parameters.
 - **Parameter number**: we collect more than the cryptographic parameters.

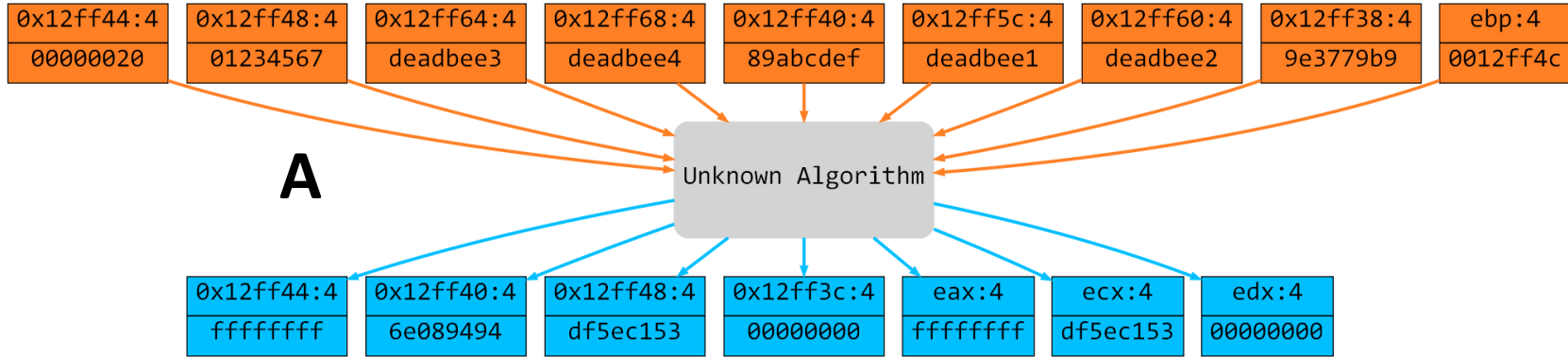
Brute-Force!





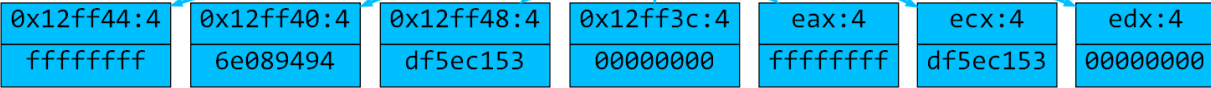
1. Generate all possible values with **A** input parameters:

1. Length 4: *00000020, 01234567, deadbee3...*
2. Length 8: *0000002001234567, 00000020deadbee3,..*
3. ...



A

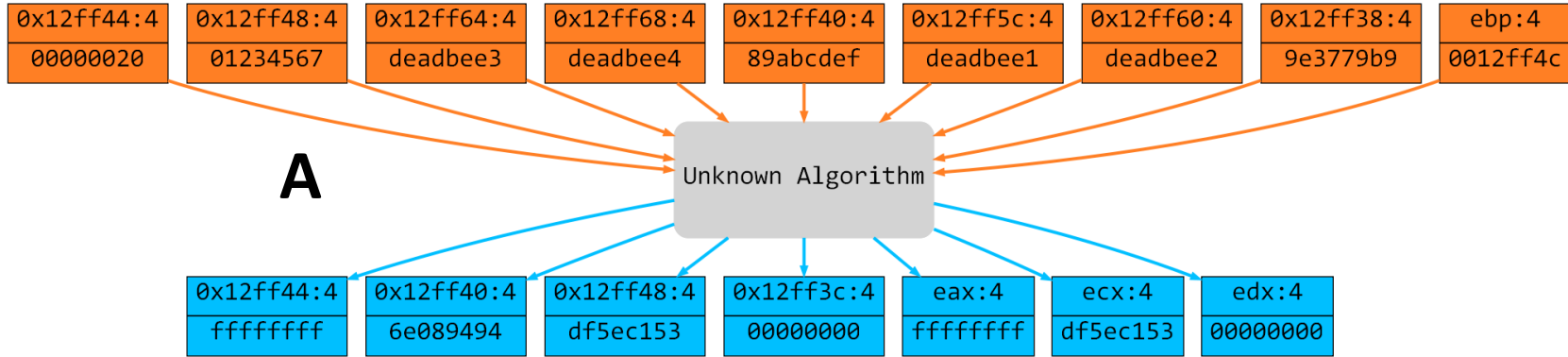
Unknown Algorithm



1. Generate all possible values with **A** input parameters:

1. Length 4: *00000020, 01234567, deadbee3...*
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3. ...

2. Same thing with **A** output parameters.



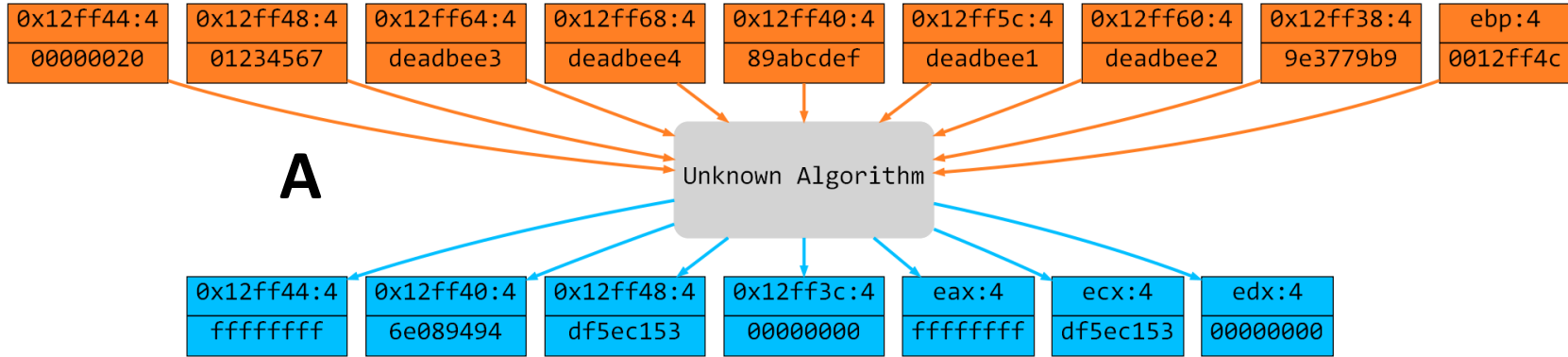
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3. For TEA reference implementation:

1. Possible input texts (8 bytes): *0000002001234567,..*



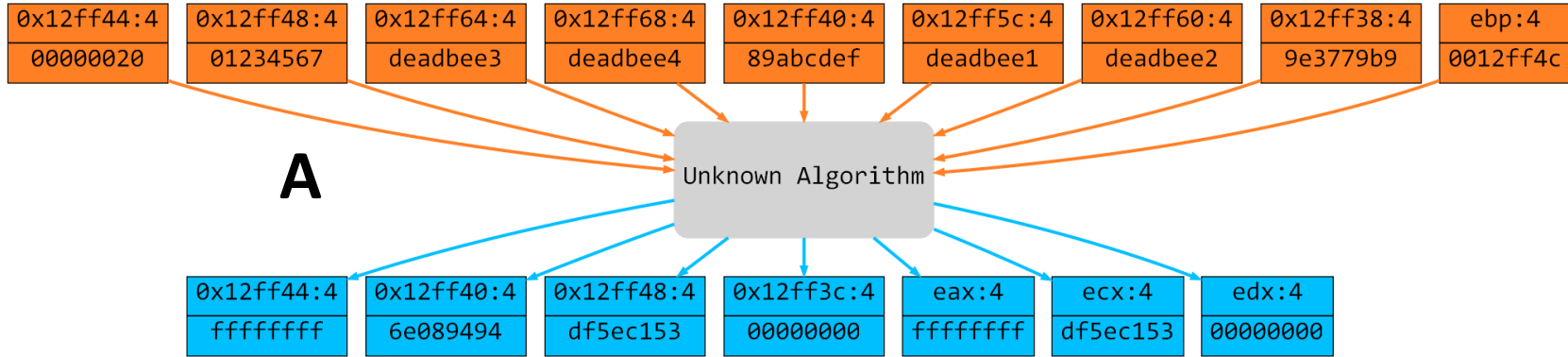
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3. For TEA reference implementation:

1. Possible input texts (8 bytes): *0000002001234567,..*
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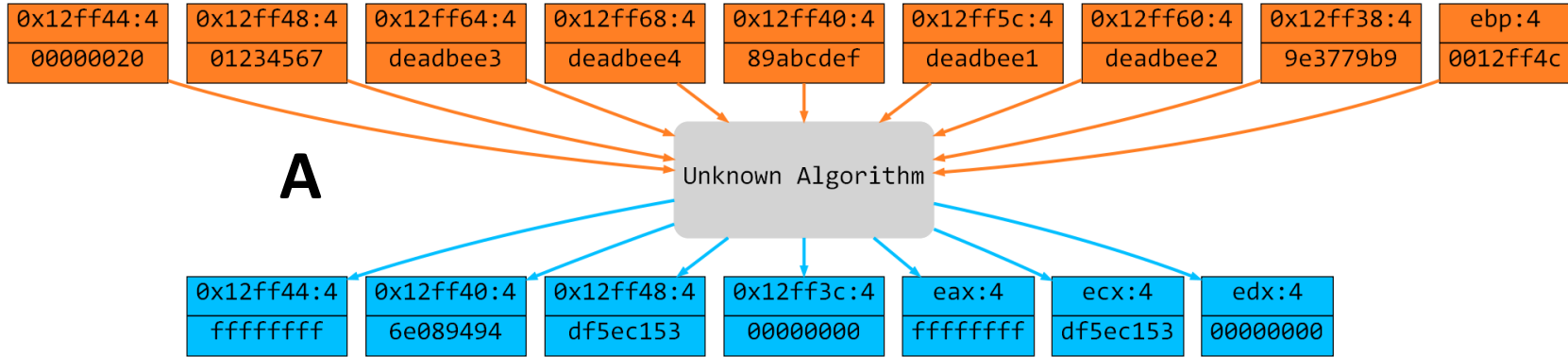
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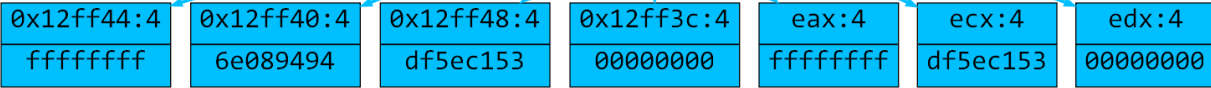
3. For TEA reference implementation:

1. Possible input texts (8 bytes): *0000002001234567,..*
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3. Execute our TEA reference implementation on each possible pair (input text, key)



A

Unknown Algorithm



1. Generate all possible values with **A** input parameters:

1. Length 4: *00000020, 01234567, deadbee3...*
2. Length 8: *0000002001234567, 00000020deadbee3,..*
3. ...

2. Same thing with **A** output parameters.

3. For TEA reference implementation:

1. Possible input texts (8 bytes): *0000002001234567,..*
2. Possible keys (16 bytes): ...
3. Execute our TEA reference implementation on each possible pair (input text, key)
4. If the output has been produced during step 2: success!


```
STARTED AT
2012-04-08 08:59:58.284000
** Crypto Algorithm Identification starting !**
9 input parameters
7 output parameters
All possible input values generation... Done!
All possible output values generation... Done!
Build internal structure... Done!
Comparison phase starting... Test for TEA decryption...
** Found TEA decryption **

==> Key (16 bytes) : deadbee1deadbee2deadbee3deadbee4

==> Crypted text (8 bytes) : 0123456789abcdef

==> Decrypted text (8 bytes) : df5ec1536e089494

ENDED AT
2012-04-08 09:01:37.832000
```

~ 2 minutes

Malware And TEA

EXAMPLES!

Storm Worm

- Several internet references about the use of TEA in the Storm Worm packer (aka Tibs).
- Let's take a look to the code...

```
push    edi
mov     ebx, [edi]
mov     ecx, [edi+4]
mov     edx, 9E3779B9h
mov     eax, edx
shl     eax, 5
mov     edi, 20h ; ' '
```

```
loc_1FA1:
mov     ebp, ebx
shl     ebp, 4
sub     ecx, ebp
mov     ebp, [esi+8]
xor     ebp, ebx
sub     ecx, ebp
mov     ebp, ebx
shr     ebp, 5
xor     ebp, eax
sub     ecx, ebp
sub     ecx, [esi+0Ch]
mov     ebp, ecx
shl     ebp, 4
sub     ebx, ebp
mov     ebp, [esi]
xor     ebp, ecx
sub     ebx, ebp
mov     ebp, ecx
shr     ebp, 5
xor     ebp, eax
sub     ebx, ebp
sub     ebx, [esi+4]
sub     eax, edx
dec     edi
jnz     short loc_1FA1
```

```
pop     edi
mov     [edi], ebx
mov     [edi+4], ecx
retn
```

```
push    edi
mov     ebx, [edi]
mov     ecx, [edi+4]
mov     edx, 9E3779B9h
mov     eax, edx
shl    eax, 5
mov     edi, 20h ;
```

TEA delta

TEA round number

```
loc_1FA1:
mov     ebp, ebx
shl    ebp, 4
sub     ecx, ebp
mov     ebp, [esi+8]
xor     ebp, ebx
sub     ecx, ebp
mov     ebp, ebx
shr    ebp, 5
xor     ebp, eax
sub     ecx, ebp
sub     ecx, [esi+0Ch]
mov     ebp, ecx
shl    ebp, 4
sub     ebx, ebp
mov     ebp, [esi]
xor     ebp, ecx
sub     ebx, ebp
mov     ebp, ecx
shr    ebp, 5
xor     ebp, eax
sub     ebx, ebp
sub     ebx, [esi+4]
sub     eax, edx
dec     edi
jnz    short loc_1FA1
```

```
pop     edi
mov     [edi], ebx
mov     [edi+4], ecx
retn
```

Classic TEA operations

```
push    edi
mov     ebx, [edi]
mov     ecx, [edi+4]
mov     edx, 9E3779B9h
mov     eax, edx
shl     eax, 5
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```

TEA delta

TEA round number

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loc_1FA1:
mov     ebp, ebx
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mov     ebp, ebx
shr     ebp, 5
xor     ebp, eax
sub     ecx, ebp
sub     ecx, [esi+0Ch]
mov     ebp, ecx
shl     ebp, 4
sub     ebx, ebp
mov     ebp, [esi]
xor     ebp, ecx
sub     ebx, ebp
mov     ebp, ecx
shr     ebp, 5
xor     ebp, eax
sub     ebx, ebp
sub     ebx, [esi+4]
sub     eax, edx
dec     edi
jnz     short loc_1FA1
```

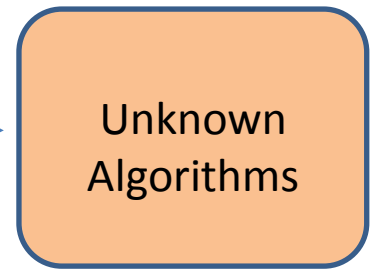
Classic TEA operations

Let's try our tool...

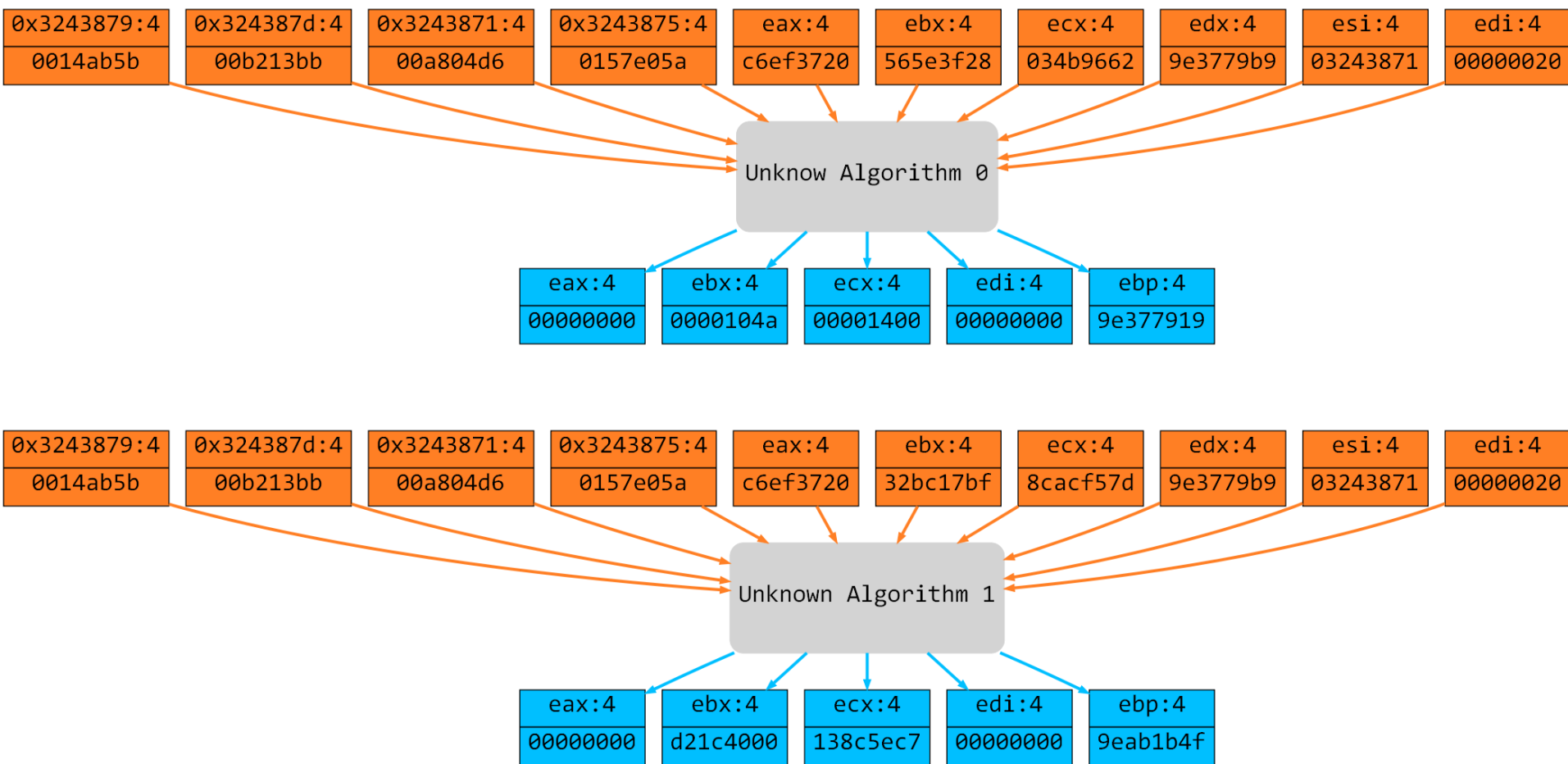
```
pop     edi
mov     [edi], ebx
mov     [edi+4], ecx
retn
```



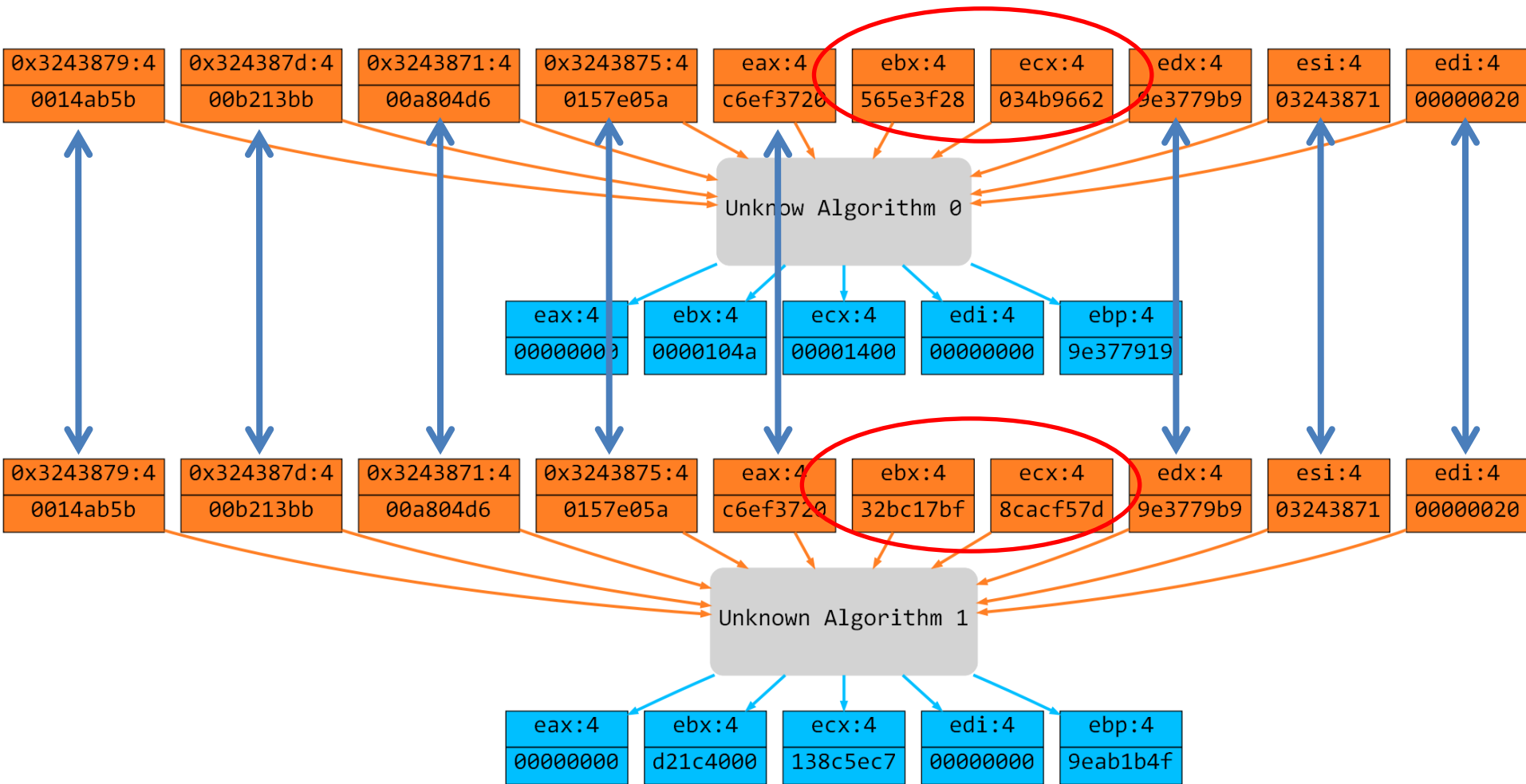
Storm
Worm
Sample



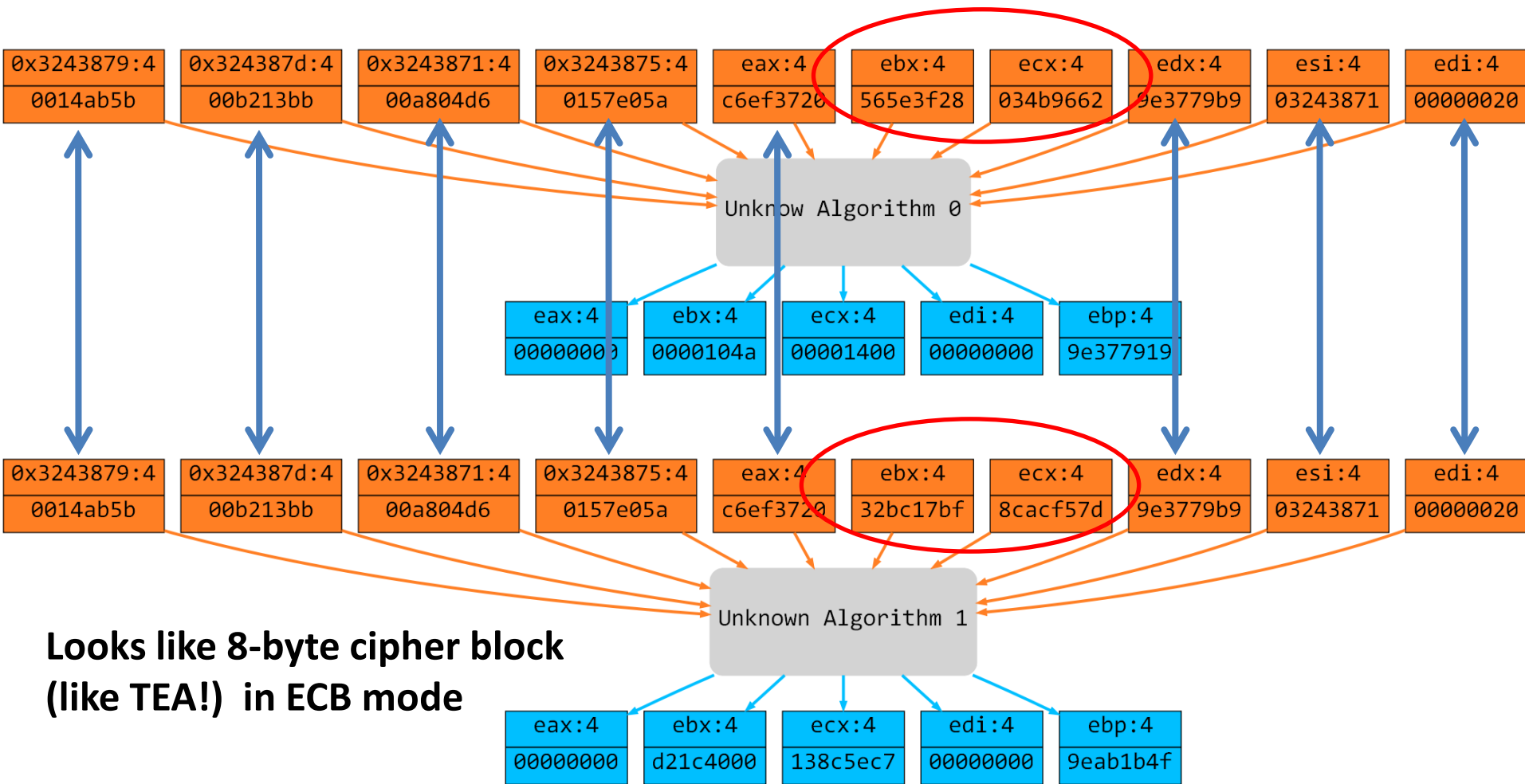
- For the previous loop, we extracted many unknown algorithms like these ones:



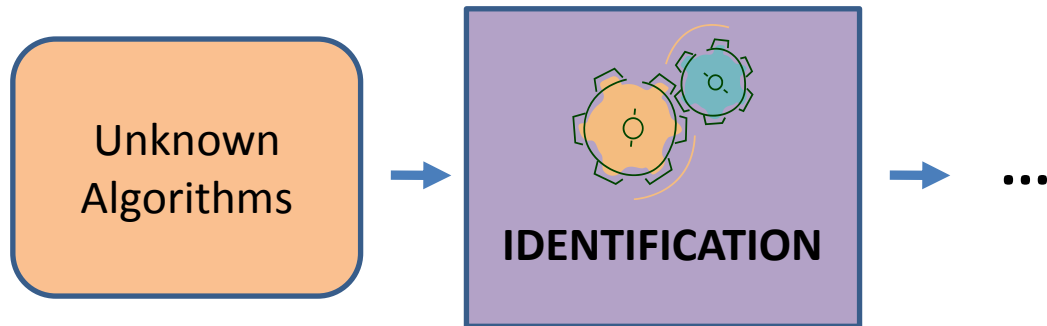
- For the previous loop, we extracted many unknown algorithms like these ones:



- For the previous loop, we extracted many unknown algorithms like these ones:



**Looks like 8-byte cipher block
(like TEA!) in ECB mode**



```
STARTED AT
2012-04-08 14:20:30.858000
** Crypto Algorithm Identification starting !**
8 input parameters
5 output parameters
All possible input values generation... Done!
All possible output values generation... Done!
Build internal structure... Done!
Comparison phase starting... Test for TEA decryption... Unknown Algorithm stays
unkown!
Done!
ENDED AT
2012-04-08 14:21:11.328000
```

WTF ?

Original TEA source code

```
z -= ((y << 4) + k[2]) ^ (y + sum) ^ ((y >> 5) + k[3])
```

```
z -= 16 * y + (y ^ *(key + 8)) + (sum ^ (y >> 5)) + *(key + 12)
```

Storm Worm implementation

Original TEA source code

```
z -= ((y << 4) + k[2]) ^ (y + sum) ^ ((y >> 5) + k[3])
```

```
z -= 16 * y + (y ^ *(key + 8)) + (sum ^ (y >> 5)) + *(key + 12)
```

Storm Worm implementation

Original TEA source code

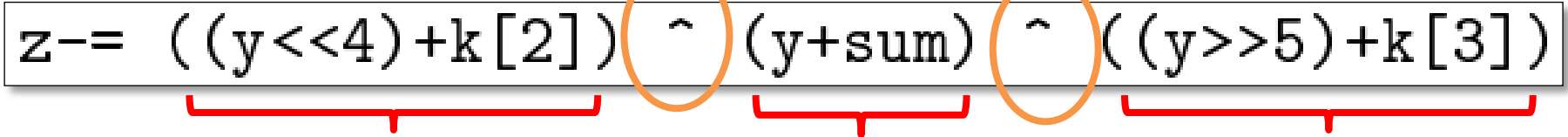
```
z -= ((y << 4) + k[2]) ^ (y + sum) ^ ((y >> 5) + k[3])
```

```
z -= 16 * y + (y ^ *(key + 8)) + (sum ^ (y >> 5)) + *(key + 12)
```

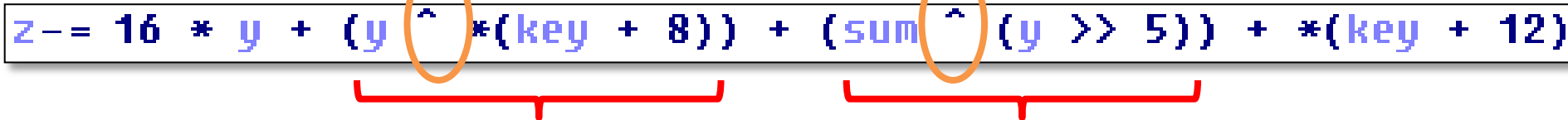
Storm Worm implementation

Original TEA source code

```
z -= ((y << 4) + k[2]) ^ (y + sum) ^ ((y >> 5) + k[3])
```

The original TEA source code line is shown in a white box with a black border. The expression is `z -= ((y << 4) + k[2]) ^ (y + sum) ^ ((y >> 5) + k[3])`. Three red brackets are positioned below the code, each under one of the three terms being XORed: `((y << 4) + k[2])`, `(y + sum)`, and `((y >> 5) + k[3])`. The two XOR operators (`^`) are circled in orange.

```
z -= 16 * y + (y ^ *(key + 8)) + (sum ^ (y >> 5)) + *(key + 12)
```

The Storm Worm implementation line is shown in a white box with a black border. The expression is `z -= 16 * y + (y ^ *(key + 8)) + (sum ^ (y >> 5)) + *(key + 12)`. Two red brackets are positioned below the code, each under one of the XORed terms: `(y ^ *(key + 8))` and `(sum ^ (y >> 5))`. The two XOR operators (`^`) are circled in orange.

Storm Worm implementation

This is not TEA: parenthesis at the wrong place!


```
STARTED AT
2012-04-10 13:22:02.192000
** Crypto Algorithm Identification starting !**
8 input parameters
5 output parameters
All possible input values generation... Done!
All possible output values generation... Done!
Build internal structure... Done!
Comparison phase starting... Test for *Storm Worm* TEA decryption...
** Found *Storm Worm* TEA decryption **

==> Key (16 bytes) : 00a804d60157e05a0014ab5b00b213bb
==> Crypted text (8 bytes) : 565e3f28034b9662
==> Decrypted text (8 bytes) : 0000104a00001400
ENDED AT
2012-04-10 13:22:13.877000
```



**Ok, Storm Worm implementation added to the base...
(this is *not* TEA)**

Trojan.SilentBanker

- Several internet references about the use of TEA in SilentBanker.
- Let's take a look to the code...

(sounds familiar, isn't it ?)

```
mov     [ebp+arg_0], 0C6EF3720h
mov     [ebp+var_C], edi
mov     [ebp+var_10], eax
mov     [ebp+arg_4], 20h
```

```
loc_4012EC:
mov     edi, [ebp+var_4]
mov     eax, edx
shr     eax, 5
xor     eax, [ebp+arg_0]
xor     edi, edx
add     edi, [ebp+var_8]
mov     ebx, edx
shl     ebx, 4
add     edi, eax
add     ebx, edi
mov     edi, [ebp+var_C]
sub     ecx, ebx
mov     eax, ecx
xor     edi, ecx
add     edi, [ebp+var_10]
mov     ebx, ecx
shr     eax, 5
xor     eax, [ebp+arg_0]
add     [ebp+arg_0], 61C88647h
shl     ebx, 4
add     edi, eax
add     ebx, edi
sub     edx, ebx
dec     [ebp+arg_4]
jnz     short loc_4012EC
```

```
mov     [esi], edx
mov     [esi+4], ecx
pop     edi
pop     esi
pop     ebx
```

```
mov [ebp+arg_0], 0C6EF3720h
mov [ebp+var_C], edi
mov [ebp+var_10], eax
mov [ebp+arg_4], 20h
```

TEA classic constant
($\delta * \text{round number}$)

TEA round number

```
loc_4012EC:
mov edi, [ebp+var_4]
mov eax, edx
shr eax, 5
xor eax, [ebp+arg_0]
xor edi, edx
add edi, [ebp+var_8]
mov ebx, edx
shl ebx, 4
add edi, eax
add ebx, edi
mov edi, [ebp+var_C]
sub ecx, ebx
mov eax, ecx
xor edi, ecx
add edi, [ebp+var_10]
mov ebx, ecx
shr eax, 5
xor eax, [ebp+arg_0]
add [ebp+arg_0], 61C88647h
shl ebx, 4
add edi, eax
add ebx, edi
sub edx, ebx
dec [ebp+arg_4]
jnz short loc_4012EC
```

= sub [ebp+arg_0], 0x9E3779B9

```
mov [esi], edx
mov [esi+4], ecx
pop edi
pop esi
pop ebx
```

Classic TEA operations

```
mov [ebp+arg_0], 0C6EF3720h
mov [ebp+var_C], edi
mov [ebp+var_10], eax
mov [ebp+arg_4], 20h
```

TEA classic constant
($\delta * \text{round number}$)

TEA round number

```
loc_4012EC:
mov edi, [ebp+var_4]
mov eax, edx
shr eax, 5
xor eax, [ebp+arg_0]
xor edi, edx
add edi, [ebp+var_8]
mov ebx, edx
shl ebx, 4
add edi, eax
add ebx, edi
mov edi, [ebp+var_C]
sub ecx, ebx
mov eax, ecx
xor edi, ecx
add edi, [ebp+var_10]
mov ebx, ecx
shr eax, 5
xor eax, [ebp+arg_0]
add [ebp+arg_0], 61C88647h
shl ebx, 4
add edi, eax
add ebx, edi
sub edx, ebx
dec [ebp+arg_4]
jnz short loc_4012EC
```

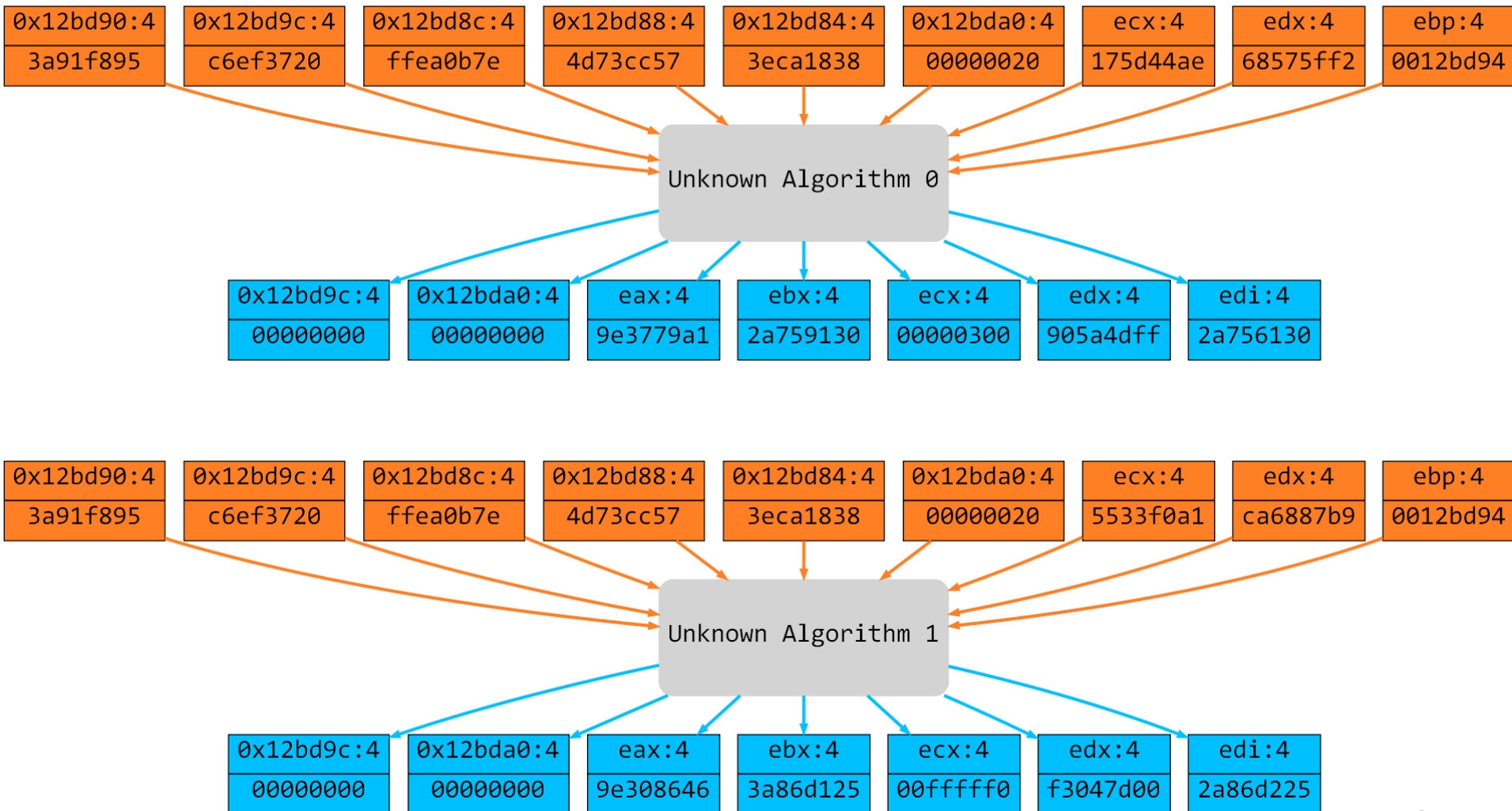
Let's try
our tool...

= sub [ebp+arg_0], 0x9E3779B9

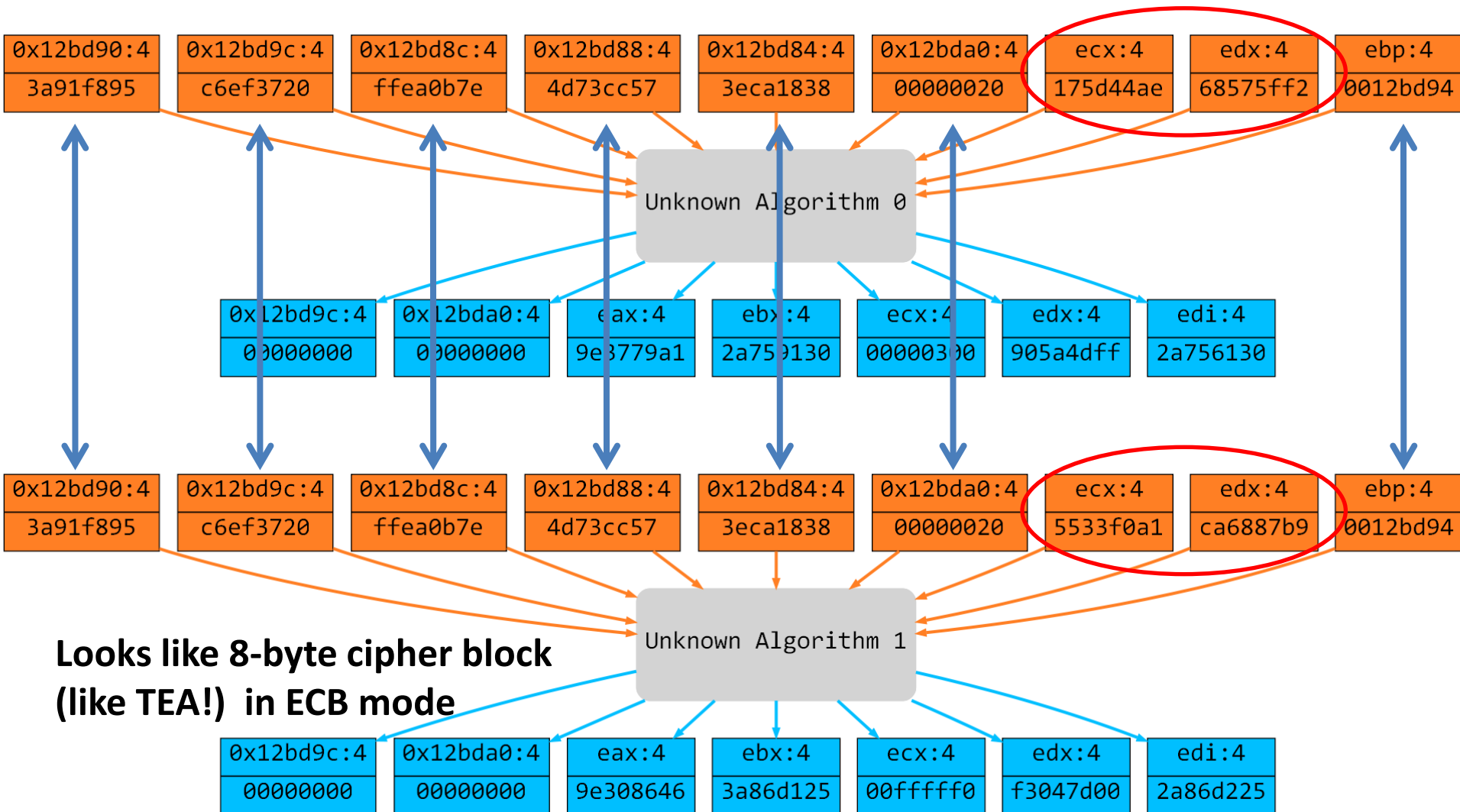
```
mov [esi], edx
mov [esi+4], ecx
pop edi
pop esi
pop ebx
```

Classic TEA operations

- For the previous loop, we extracted many unknown algorithms like these ones:



- For the previous loop, we extracted many unknown algorithms like these ones:



```
STARTED AT  
2012-04-08 16:05:14.288000  
** Crypto Algorithm Identification starting !**  
9 input parameters  
7 output parameters  
All possible input values generation... Done!  
All possible output values generation... Done!  
Build internal structure... Done!  
Comparison phase starting... Test for TEA decryption... Unknown Algorithm stays  
unknown!  
Done!
```

Fail.. Again !?


```
STARTED AT
2012-04-08 16:15:22.255000
** Crypto Algorithm Identification starting !**
9 input parameters
7 output parameters
All possible input values generation... Done!
All possible output values generation... Done!
Build internal structure... Done!
Comparison phase starting... Test for *Storm Worm* TEA decryption...
** Found *Storm Worm* TEA decryption **

==> Key (16 bytes) : 4d73cc573eca18383a91f895ffea0b7e
==> Crypted text (8 bytes) : 68575ff2175d44ae
==> Decrypted text (8 bytes) : 905a4dff300

ENDED AT
2012-04-08 16:16:22.552000
```



Same implementation than in the Storm Worm!

- They probably both copied/pasted a wrong source code from the internet.
- Started to look for it: Google, TEA Wikipedia page,... nothing!
- At some point, I remembered something these two malware families have in common...



They both came from Russia!

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Tiny Encryption Algorithm

From Wikipedia, the free encyclopedia

In [cryptography](#), the **Tiny Encryption Algorithm** (TEA) is a stream cipher consisting of a few lines of code. It was designed by [David Wheeler](#) and [Sami Shalunov](#) at the [Fast Software Encryption](#) workshop in [Leuven](#) in 1994.

The cipher is not subject to any [patents](#).

Contents [\[hide\]](#)















- [Properties](#)
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Properties

TEA operates on two 32-bit unsigned integers (could be 16-bit) and uses a suggested 64 rounds, typically implemented in pairs of exactly the same way for each cycle. Different multiple rounds. The magic constant, 2654435769 or 9E3779B5.

TEA has a few weaknesses. Most notably, it suffers from a small effective key size is only 126 bits.^[3] As a result, TEA is vulnerable to [hacking Microsoft's Xbox game console](#), where the cipher requires 2^{23} chosen plaintexts under a related-key pair designed.

Ссылки

- [Страница алгоритма шифрования TEA](#)  
- Roger M. Needham and David J. Wheeler. «TEA, a Tiny Encryption Algorithm»(PDF)  
- Andrew Hargreaves. «Hacking the Xbox: an introduction to reverse engineering»[1]  
- Журнал «Хакер Онлайн», TEA: блочный шифр своими руками, (2004)[2]  
- Paris Kitsos, Yan Zhang. «RFID Security: Techniques, Protocols and System-On-Chip Design»[3]  
- David J. Wheeler and Roger M. Needham. «Correction to xtea.» Technical report, Computer Laboratory,
- Roger M. Needham and David J. Wheeler. «Tea extensions.» Technical report, Computer Laboratory, University of Cambridge
- [Test vectors for TEA](#)  
- [JavaScript implementation of XXTEA with Base64](#)  



Взлом как образ мысли

Интервью с человеком, который, как оказалось, является не только талантливым пентестером в одной из крупных ИБ-компаний, но и хакером-ветераном, который уверенными шагами вышел на свет и прикоснулся к истории российской хак-сцены....



Полный г голосова

Конкурсы с го посетителей, интерес вызы лакомые приз

TEA: блочный шифр своими руками



Дата: **22.04.2004**    |  [slon](#)

В данном тексте хотелось бы затронуть такую животрепещущую тему, как шифрование файлов. Вообще нужно различать два вида шифрования файлов:

- шифрование для себя (чтобы ваши файлы никто, кроме вас не "читал")
- шифрование для других (чтобы ваши файлы "читал" только адресат)



Оценка:

 **11** 

статьи

 Tweet

Russian Website TEA source code

```
push edi ; Сохраняем  
mov ebx,v0 ; Кладём  
mov ecx,v1 ; В ecx кл  
mov edx,9e3779b9h ;  
mov eax,edx ; Кладём  
shl eax,5 ; Сдвиг eax  
mov edi,32 ; Кладём  
DLoopR:  
mov ebp,ebx ; Кладём  
shl ebp,4 ; Сдвиг ebp  
sub ecx,ebp ; Отнима  
mov ebp,k2 ; Кладём  
xor ebp,ebx ; XOR'им  
sub ecx,ebp ; Отнима  
mov ebp,ebx ; Кладём  
shr ebp,5 ; Сдвиг ebp  
xor ebp,eax ; XOR'им  
sub ecx,ebp ; Отнима  
sub ecx,k3 ; Отнимаем  
;  
mov ebp,ecx ; Кладём  
shl ebp,4 ; Сдвиг ebp  
sub ebx,ebp ; Отнима  
mov ebp,k0 ; Кладём  
xor ebp,ecx ; XOR'им  
sub ebx,ebp ; Отнима  
mov ebp,ecx ; Кладём  
shr ebp,5 ; Сдвиг ebp  
xor ebp,eax ; XOR'им  
sub ebx,ebp ; Отнима  
sub ebx,k1 ; Отнимаем  
sub eax,edx ; Отнима  
dec edi ; Уменьшаем  
jnz DLoopR ; Дешифр  
  
pop edi ; Вынимаем  
mov v0,ebx ; Кладём  
mov v1,ecx ; В отведё  
ret ; Возврат из подп
```

Russian
Website
TEA source
code

```
push edi ; Сохраняем и  
mov ebx,v0 ; Кладём  
mov ecx,v1 ; В ecx кл  
mov edx,9e3779b9h ;  
mov eax,edx ; Кладём  
shl eax,5 ; Сдвиг eax  
mov edi,32 ; Кладём и  
DLoopR:  
mov ebp,ebx ; Кладём  
shl ebp,4 ; Сдвиг ebp  
sub ecx,ebp ; Отнима  
mov ebp,k2 ; Кладём  
xor ebp,ebx ; XOR'им  
sub ecx,ebp ; Отнима  
mov ebp,ebx ; Кладём  
shr ebp,5 ; Сдвиг ebp  
xor ebp,eax ; XOR'им  
sub ecx,ebp ; Отнима  
sub ecx,k3 ; Отнимаем  
;  
mov ebp,ecx ; Кладём  
shl ebp,4 ; Сдвиг ebp  
sub ebx,ebp ; Отнима  
mov ebp,k0 ; Кладём  
xor ebp,ecx ; XOR'им  
sub ebx,ebp ; Отнима  
mov ebp,ecx ; Кладём  
shr ebp,5 ; Сдвиг ebp  
xor ebp,eax ; XOR'им  
sub ebx,ebp ; Отнима  
sub ebx,k1 ; Отнимаем  
sub eax,edx ; Отнима  
dec edi ; Уменьшаем  
jnz DLoopR ; Дешифр  
  
pop edi ; Вынимаем и  
mov v0,ebx ; Кладём  
mov v1,ecx ; В отведё  
ret ; Возврат из подп
```

==

```
push edi  
mov ebx, [edi]  
mov ecx, [edi+4]  
mov edx, 9E3779B9h  
mov eax, edx  
shl eax, 5  
mov edi, 20h ; ' '
```

```
loc_1FA1:  
mov ebp, ebx  
shl ebp, 4  
sub ecx, ebp  
mov ebp, [esi+8]  
xor ebp, ebx  
sub ecx, ebp  
mov ebp, ebx  
shr ebp, 5  
xor ebp, eax  
sub ecx, ebp  
sub ecx, [esi+0Ch]  
mov ebp, ecx  
shl ebp, 4  
sub ebx, ebp  
mov ebp, [esi]  
xor ebp, ecx  
sub ebx, ebp  
mov ebp, ecx  
shr ebp, 5  
xor ebp, eax  
sub ebx, ebp  
sub ebx, [esi+4]  
sub eax, edx  
dec edi  
jnz short loc_1FA1
```

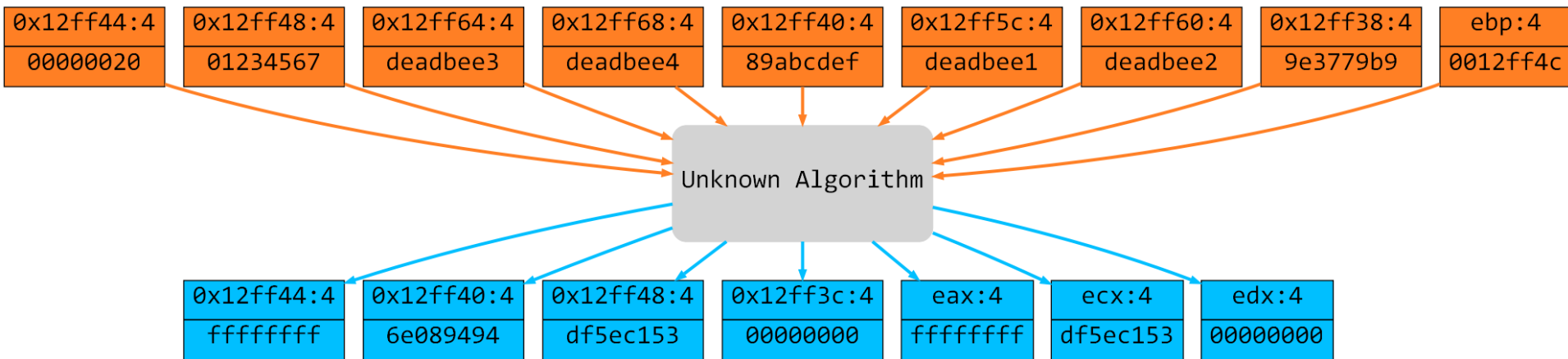
```
pop edi  
mov [edi], ebx  
mov [edi+4], ecx  
retn
```

Storm
Worm

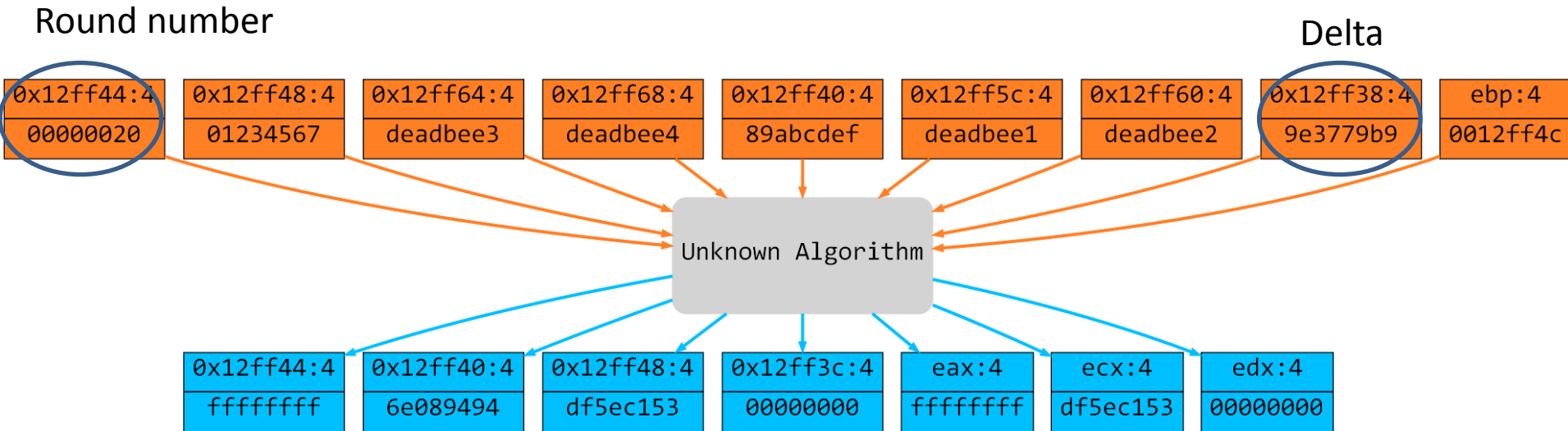
Modified TEA

MORE EXAMPLES!

Remember This ?



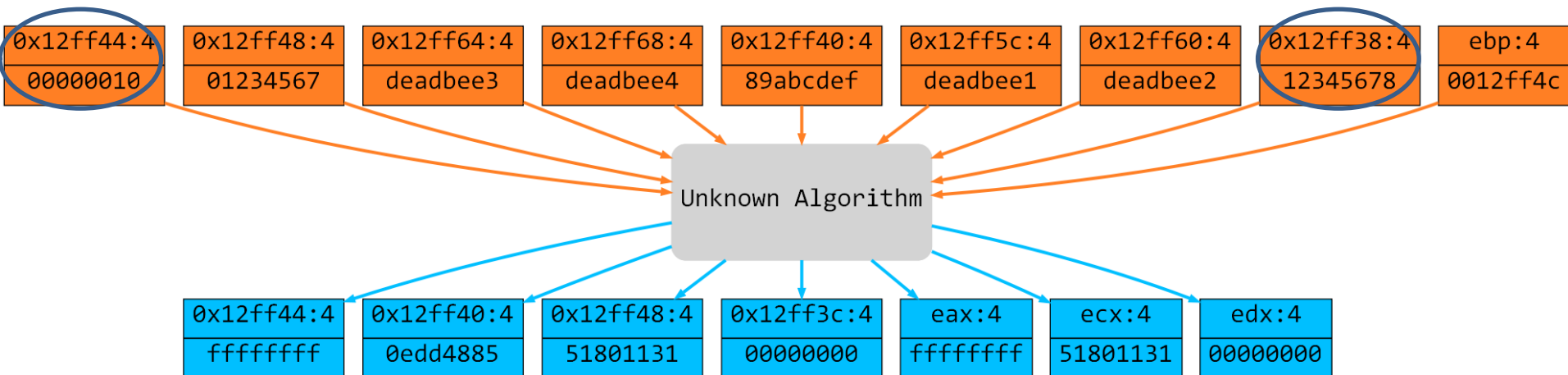
Remember This ?



The magic TEA constant (*delta*) and the round number are seen as input parameters, because they are initialized *before* the loop and used inside.

Modified TEA Implementation

- delta = 0x12345678 (normally 0x9E3779B9)
- round number = 16 (normally 32)



- TEA reference implementation extended:

```
def tea (input_text, key):
```

```
    ...
```

- TEA reference implementation extended:

```
def tea (input_text, key, delta, round_number):  
    ...
```

- TEA reference implementation extended:

```
def tea (input_text, key, delta, round_number):  
    ...
```

```
STARTED AT  
2012-04-08 13:16:27.918000  
** Crypto Algorithm Identification starting !**  
8 input parameters  
3 output parameters  
All possible input values generation... Done!  
All possible output values generation... Done!  
Build internal structure... Done!  
Comparison phase starting... Test for TEA *modified* decryption  
** Found *modified* TEA decryption **  
  
==> Key (16 bytes) :   deadbee1deadbee2deadbee3deadbee4  
  
==> Delta (4 bytes) :   12345678  
  
==> Round number (4 bytes) :   00000010  
  
==> Crypted text (8 bytes) :   0123456789abcdef  
  
==> Decrypted text (8 bytes) : 51801131edd4885  
ENDED AT  
2012-04-08 13:21:11.562000
```

RC4

MORE EXAMPLES!

RC4 (1)

- RC4 algorithm:
 - Stream cipher
 - Variable-length key
 - Two loops generate a pseudorandom stream into a 256 bytes substitution-box (S-BOX).
 - A final loop does the actual decryption.
- We have to extend our model to **regroup different loops into a same algorithm.**

Interlude: Loop Data-Flow

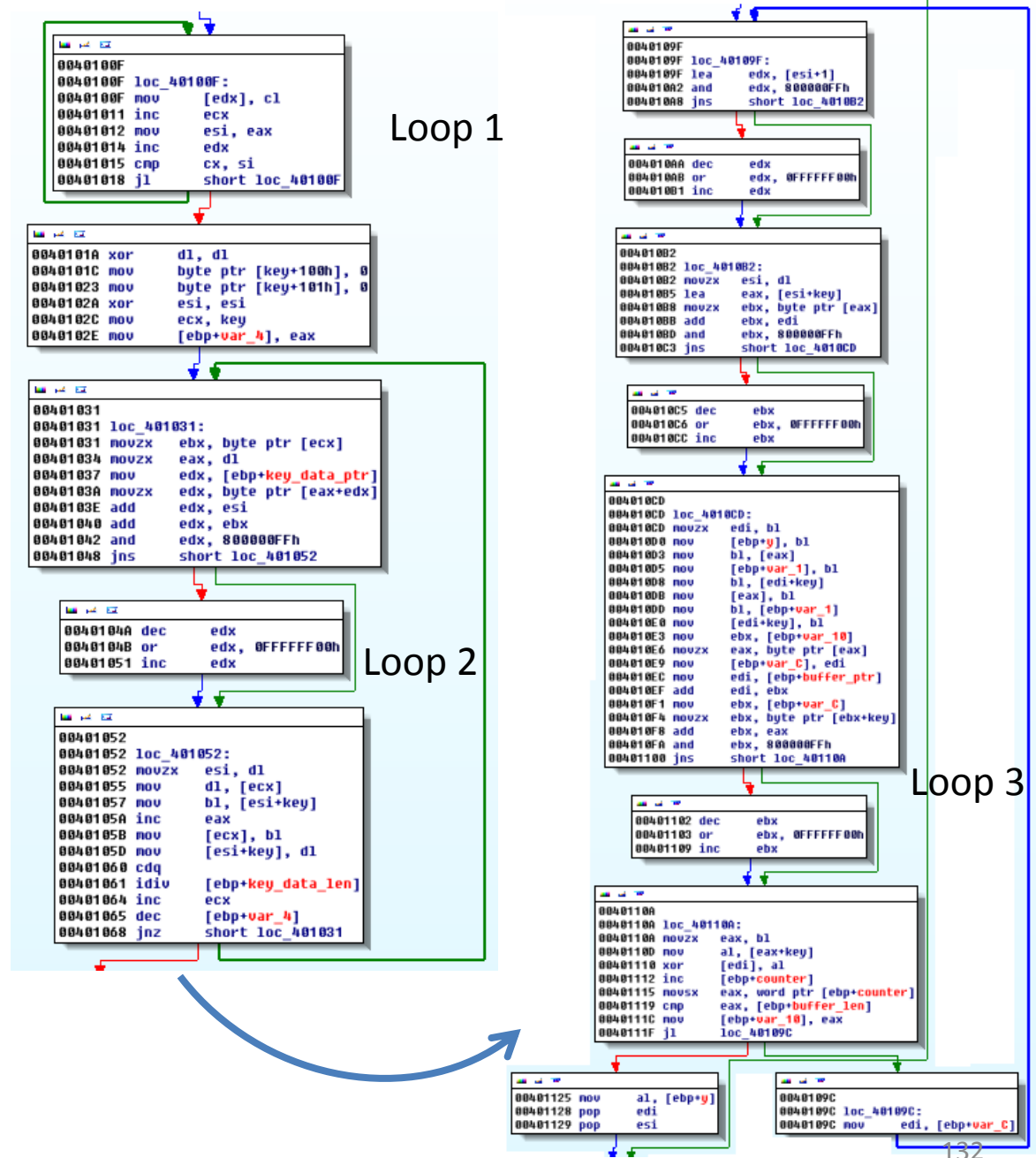
- Two loops $L1$ and $L2$ are in the same algorithm:
 - If $L1$ started before $L2$ in the trace.
 - If $L2$ uses as input parameter an output parameter of $L1$.

(or the contrary!)

RC4 (2)

- We built a toy program calling the RC4 decryption function on:
 - Key : **“SuperKeyIsASuperKey” (19 bytes)**
 - Encrypted text: **“AAA...AA” (1024 bytes)**

Statically speaking it looks like this...

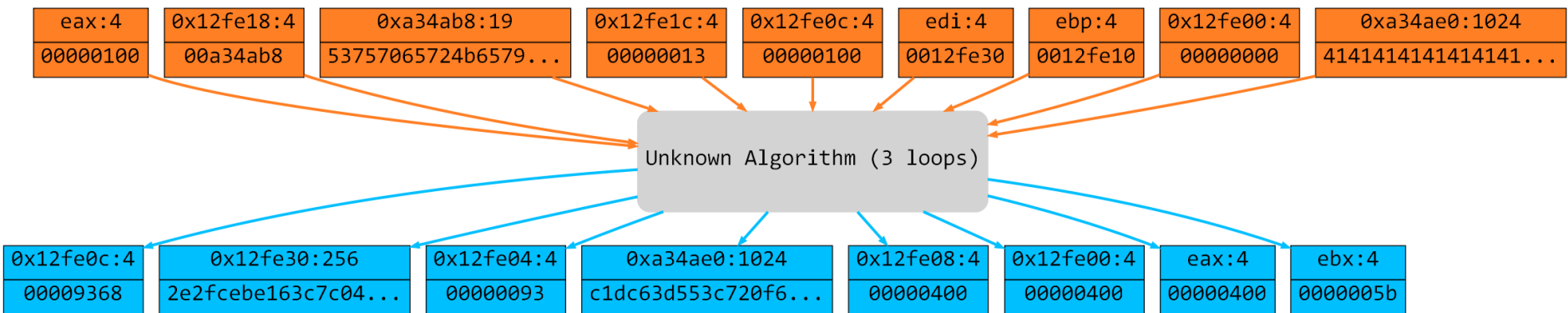


Tools	Answer
Crypto Searcher	∅
Draca v0.5.7b	∅
Findcrypt v2	∅
Hash & Crypto Detector v1.4	∅
PEiD KANAL v2.92	∅
Kerckhoffs	∅
Signsrch 0.1.7	∅
SnD Crypto Scanner v0.5b	∅

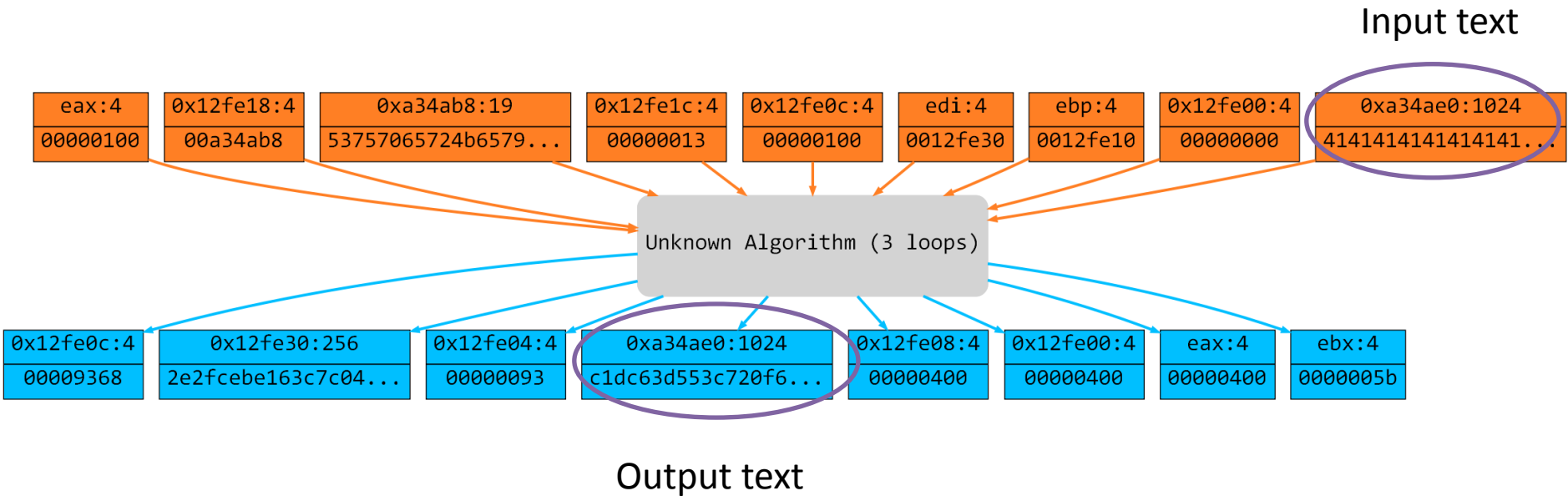
Tools	Answer
Crypto Searcher	∅
Draca v0.5.7b	∅
Findcrypt v2	∅
Hash & Crypto Detector v1.4	∅
PEiD KANAL v2.92	∅
Kerckhoffs	∅
Signsrch 0.1.7	∅
SnD Crypto Scanner v0.5b	∅

Let's try our tool...

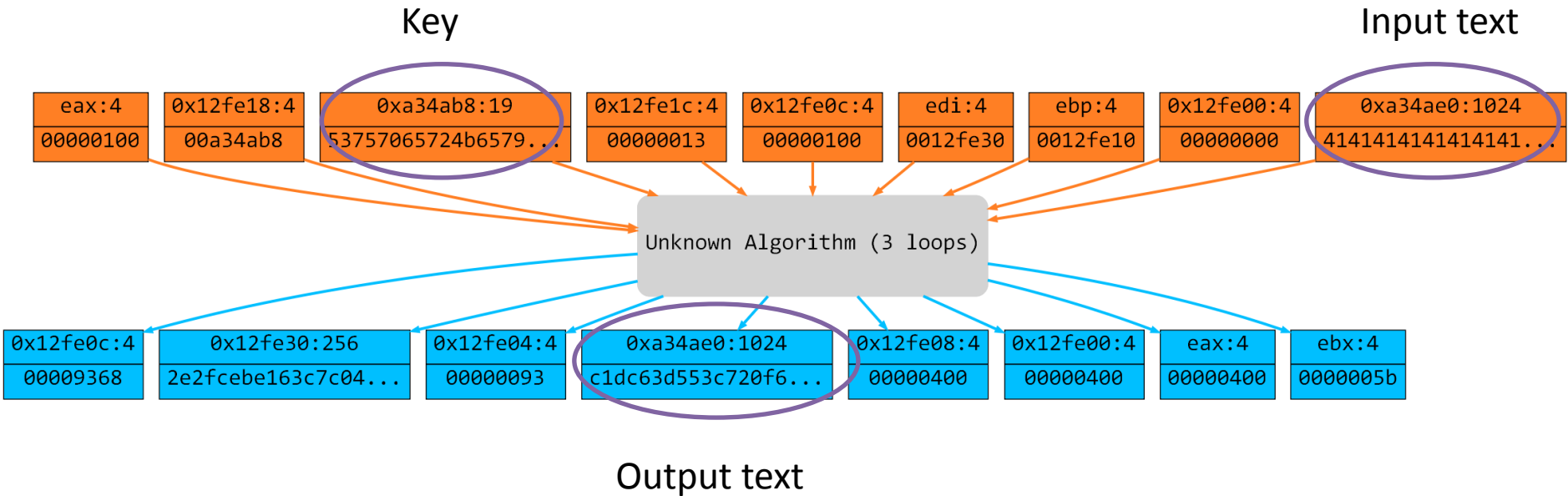
- We extracted from our toy program:



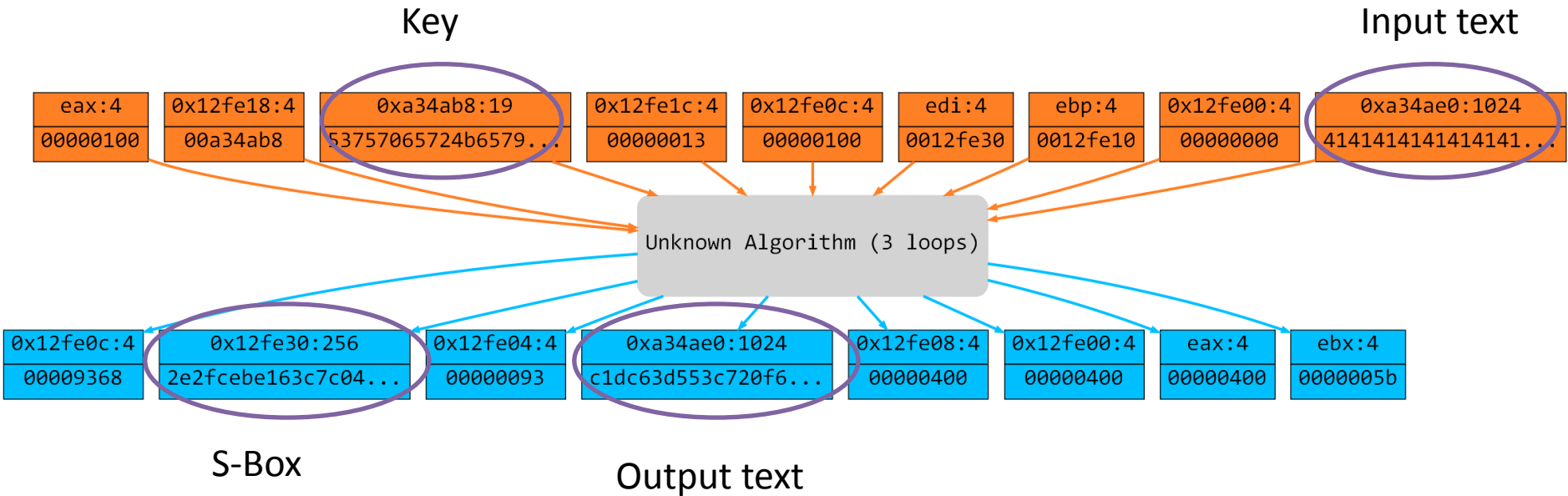
- We extracted from our toy program:



- We extracted from our toy program:



- We extracted from our toy program:



```
STARTED AT
2012-04-08 19:34:12.959000
** Crypto Algorithm Identification starting !**
3 input parameters
3 output parameters
All possible input values generation... Done!
All possible output values generation... Done!
Build internal structure... Done!
Comparison phase starting... Test for RC4...
```

```
** Found RC4 **
```

```
==> Key (19 bytes) : 53757065724b657949734153757065724b6579
```

```
==> Crypted text (1024 bytes) : 4141414141414141...
```

```
==> Decrypted text (1024 bytes) : c1dc63d553c720f6...
```

```
ENDED AT
2012-04-08 19:34:12.990000
```



Salinity

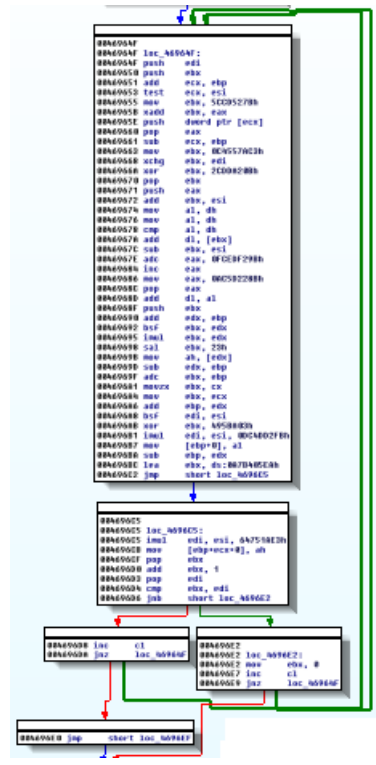
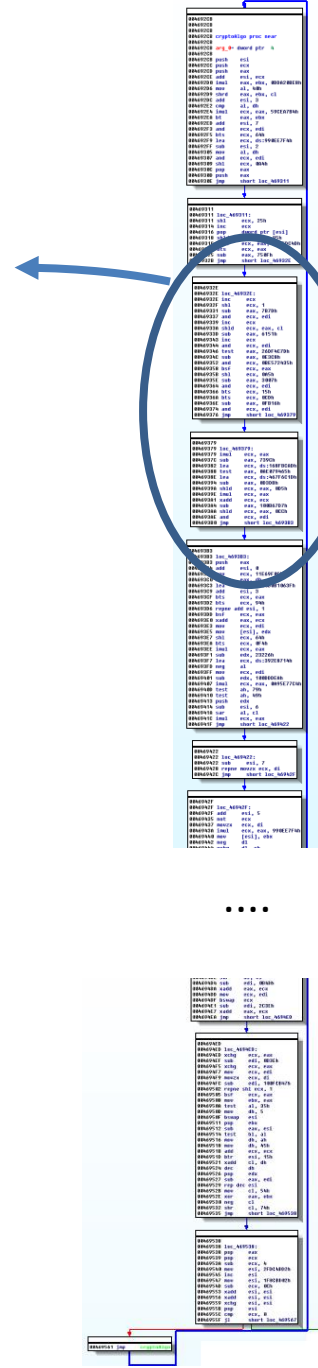
- Several internet references about the use of RC4 in the Salinity packer.
- Let's take a look...

(suspense...)

Loop 1

```
0046932E
0046932E loc_46932E:
0046932E inc     ecx
0046932F shl     ecx, 1
00469331 sub     eax, 7B7Dh
00469337 and     ecx, edi
00469339 inc     ecx
0046933A shld   ecx, eax, cl
0046933D sub     eax, 6151h
00469343 inc     ecx
00469344 and     ecx, edi
00469346 test   eax, 26DF4C7Dh
0046934C sub     eax, 0E3C8h
00469352 and     ecx, 0BE572435h
00469358 bsf     ecx, eax
0046935B shl     ecx, 0A5h
0046935E sub     eax, 3007h
00469364 and     ecx, edi
00469366 bts     ecx, 15h
0046936A bts     ecx, 0CDh
0046936E sub     eax, 0FB16h
00469374 and     ecx, edi
00469376 jmp     short loc_469379
```

```
00469379
00469379 loc_469379:
00469379 imul   ecx, eax
0046937C sub     eax, 739Ch
00469382 lea    ecx, ds:168FBCADh
00469388 test   eax, 0AE079465h
0046938E lea    ecx, ds:467F6C1Dh
00469394 sub     eax, 0D3DBh
0046939A shld   ecx, eax, 0D5h
0046939E imul   ecx, eax
004693A1 xadd   ecx, ecx
004693A4 sub     eax, 100867D7h
004693AA shld   ecx, eax, 0CCh
004693AE and     ecx, edi
004693B0 jmp     short loc_4693B3
```



Loop 2

```
00469722
00469722 loc_469722:
00469722 inc     dl
00469723 push  7A7B02EEh
00469724 mov     [esp+4+arg_0], esi
00469725 add     esp, edx
00469727 neg     al
00469731 mov     al, 0Bh
00469735 add     bl, [ebp+8]
00469738 sub     esp, edx
0046973E scasd  eax, esi
00469738 mov     esi, 19E4F3Ch
00469743 mov     al, 5Bh
00469746 push  0
00469748 pop     esp, edx
0046974A test   ecx, 0B70A8550h
00469751 shld   esi, ecx, cl
00469754 mov     ecx, 19E4F3Ch
00469754 mov     al, [ebp+8]
00469759 sub     esp, edx
0046975E sub     edi, edx
0046975E shr     edi, 0B0h
00469761 xchg  edi, esi
00469763 add     ebx, ebp
00469765 mov     edi, 0BF4C78C2h
00469768 inc     esi
0046976C rop   mov     cx, [ebx]
0046976F sub     ebx, ebx
00469771 mov     edi, 0744B5F0h
00469777 dsuop
00469778 mov     edi, 0BF4C8102h
0046977F add     ebx, ebp
00469781 rcr   edi, 2Ah
00469784 movsx edi, dx
00469787 dsuop
00469788 mov     [ebx], al
0046978E sub     ebx, ebx
0046978E mov     edi, ebx
00469791 bt     edi, esi
00469792 dec     edi
00469792 add     ebx, ebp
00469794 dsuop
00469797 db   0F3h
00469797 rop   mov     [ebx], ch
0046979E sub     ebx, ebx
0046979E imul  edi, esi
004697A0 mov     edi, 572A05B0h
004697A6 bsf     esi, ecx
004697A8 add     al, ch
004697AA pop     edi
004697AC add     edi, 5
004697AF mov     al, [ebp+eax+8]
004697B3 sar     [edi], al
004697B5 sub     cl, 5
004697B8 pop     esi
004697BA dec     esi
004697BC jmp     inc_469722
```

Loop 1

Hmpf.. Let's try!

```

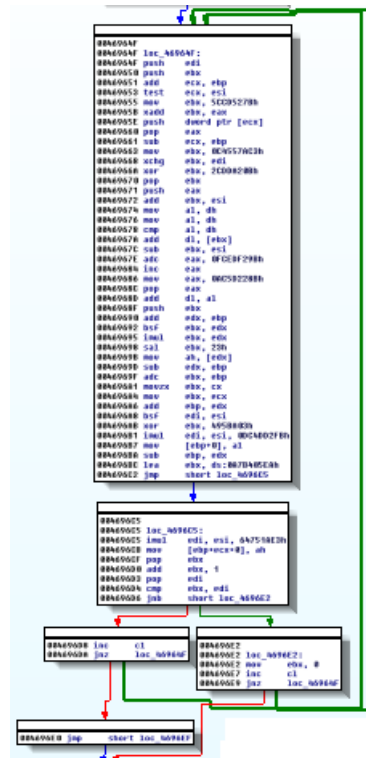
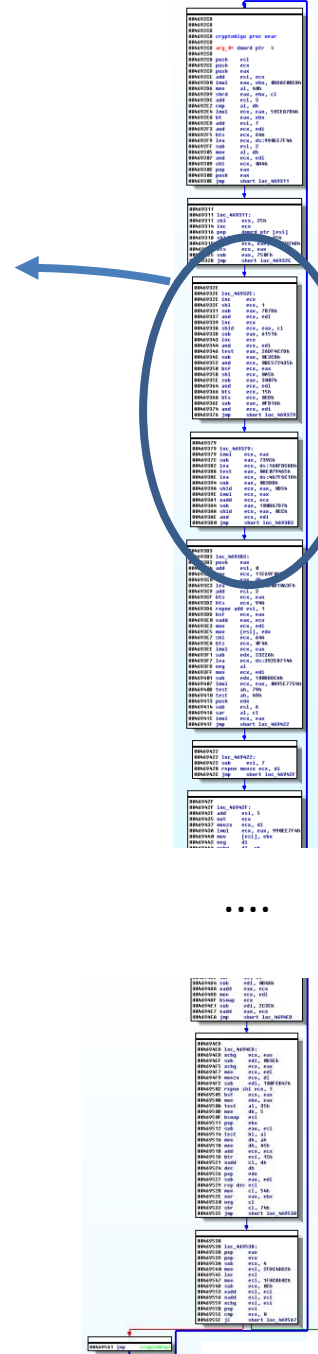
0046932E
0046932E loc_46932E:
0046932E inc     ecx
0046932F shl     ecx, 1
00469331 sub     eax, 7B7Dh
00469337 and     ecx, edi
00469339 inc     ecx
0046933A shld   ecx, eax, cl
0046933D sub     eax, 6151h
00469343 inc     ecx
00469344 and     ecx, edi
00469346 test   eax, 26DF4C7Dh
0046934C sub     eax, 0E3C8h
00469352 and     ecx, 0BE572435h
00469358 bsf     ecx, eax
0046935B shl     ecx, 0A5h
0046935E sub     eax, 3007h
00469364 and     ecx, edi
00469366 bts     ecx, 15h
0046936A bts     ecx, 0CDh
0046936E sub     eax, 0FB16h
00469374 and     ecx, edi
00469376 jmp     short loc_469379

```

```

00469379
00469379 loc_469379:
00469379 imul   ecx, eax
0046937C sub     eax, 739Ch
00469382 lea    ecx, ds:168FBCADh
00469388 test   eax, 0AE079465h
0046938E lea    ecx, ds:467F6C1Dh
00469394 sub     eax, 0D3DBh
0046939A shld   ecx, eax, 0D5h
0046939E imul   ecx, eax
004693A1 xadd   ecx, ecx
004693A4 sub     eax, 100867D7h
004693AA shld   ecx, eax, 0CCh
004693AE and     ecx, edi
004693B0 jmp     short loc_4693B3

```



Loop 3

```

00469722
00469722 loc_469722:
00469722 inc     dl
00469726 push  7A7B02EEh
00469729 mov     [esp+4+arg_0], esi
0046972B add     esp, edx
0046972F neg     al
00469731 mov     al, 0Bh
00469735 add     bl, [ebp+8]
00469738 sub     esp, edx
0046973E scasd  eax, esi
00469740 mov     esi, 19E4F3Ch
00469743 mov     al, 5Bh
00469746 push  0
00469748 pop     esp, edx
0046974B test   ecx, 0B70A8550h
00469751 shld   esi, ecx, cl
00469754 mov     ecx, 19E4F3Ch
00469756 mov     al, [ebp+8]
00469759 sub     esp, edx
0046975B sub     edi, edx
0046975E shr     edi, 0B0h
00469761 xor     edi, esi
00469763 add     ebx, ebp
00469765 mov     edi, 0BF4C70C2h
00469768 inc     esi
0046976C rop     mov     ebx, [ebx]
0046976F sub     ebx, ebp
00469771 mov     edi, 0F7A4B5F0h
00469773 dsuop  esi
00469776 mov     edi, 0BF4C5022h
00469778 add     ebx, ebp
00469781 rcr     edi, 2Ah
00469784 movsx  esi, dx
00469787 dsuop  edi
00469789 mov     [ebx], al
0046978B sub     ebx, ebp
0046978D mov     esi, ebx
0046978F bt     edi, esi
00469792 dec     edi
00469794 add     ebx, ebp
00469797 dsuop  edi
00469799 db     0F3h
0046979F rop     mov     [ebx], ch
0046979B sub     ebx, ebp
0046979E imul  edi, esi
004697A0 mov     edi, 572A05B0h
004697A6 bsf     esi, ecx
004697A8 add     al, ch
004697AA pop     edi
004697AC add     edi, 5
004697AF mov     al, [ebp+eax+8]
004697B1 sar     [edi], al
004697B5 sub     cl, 5
004697B8 pop     esi
004697BA dec     esi
004697BC jmp     loc_469722

```

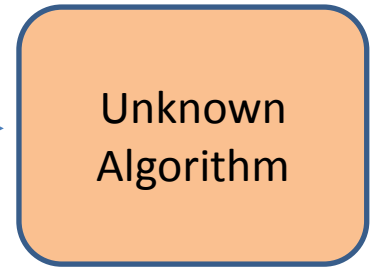
Loop 2



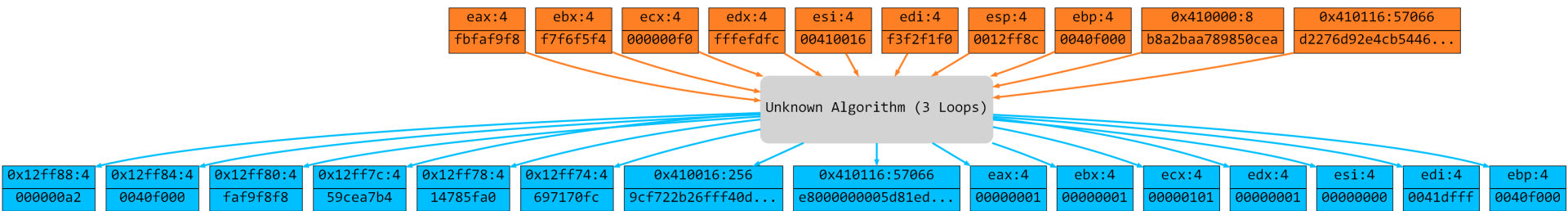
Salty
Sample



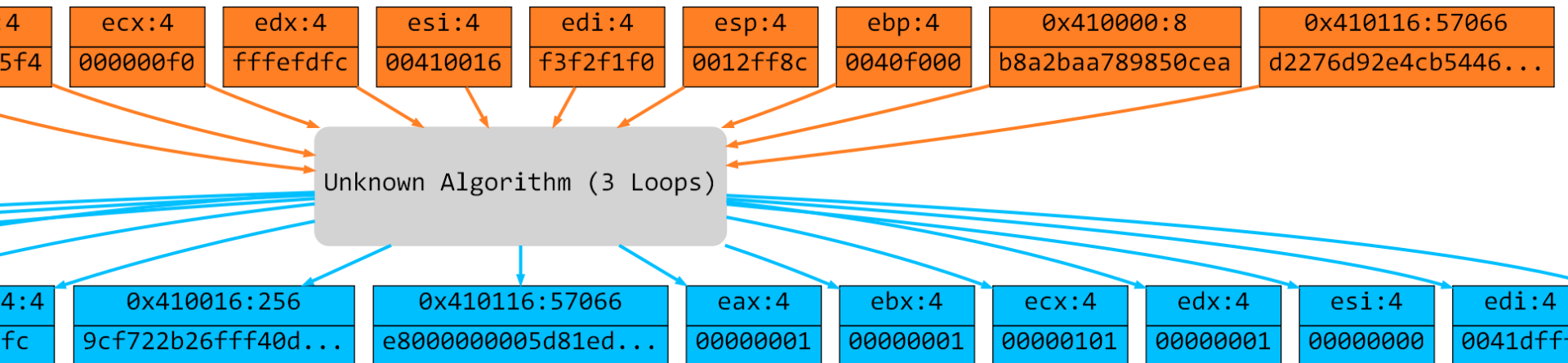
(Multi-loops)



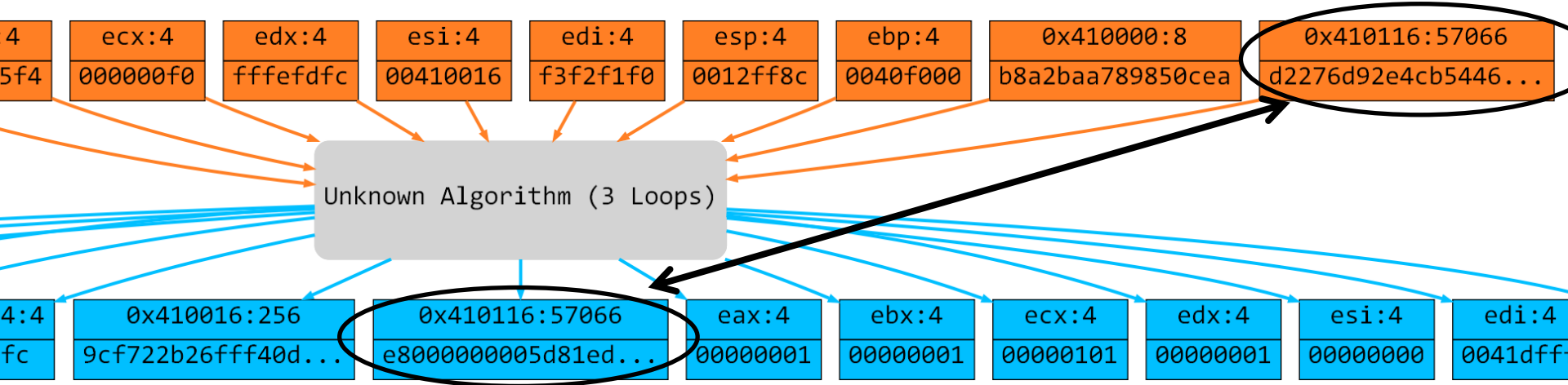
For the previous 3 loops, we extracted one algorithm:



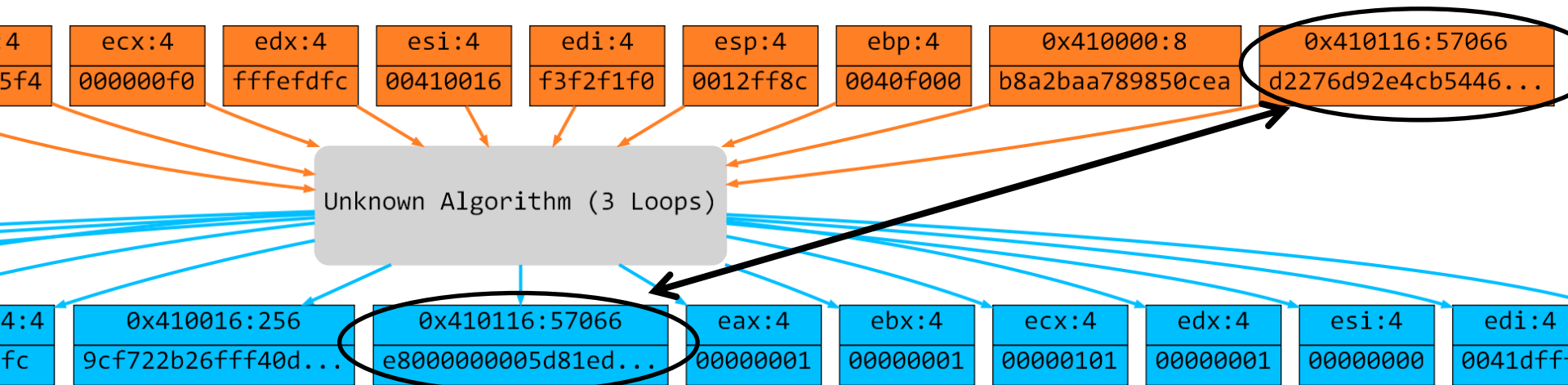
For the previous 3 loops, we extracted one algorithm:



For the previous 3 loops, we extracted one algorithm:

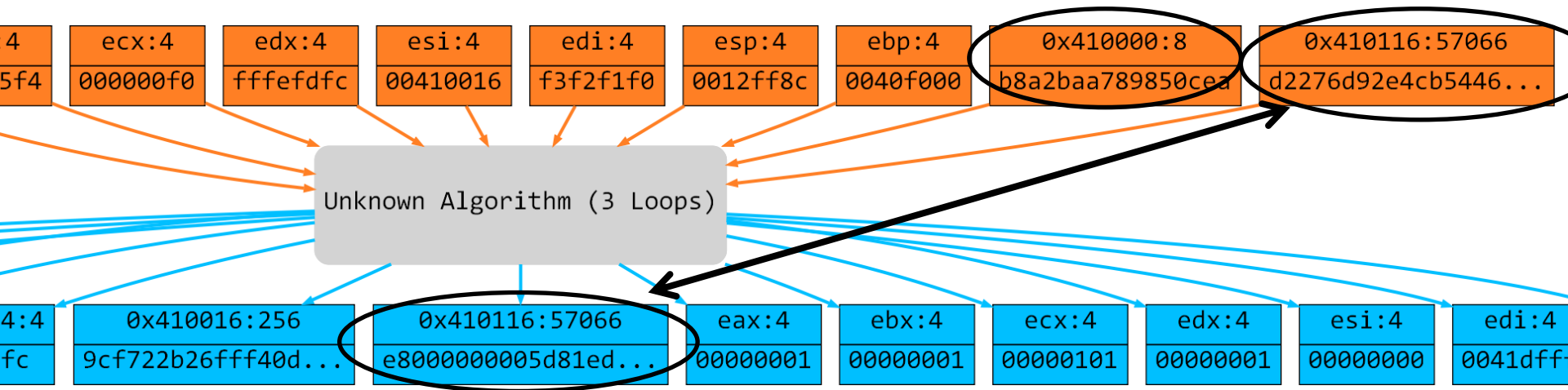


For the previous 3 loops, we extracted one algorithm:



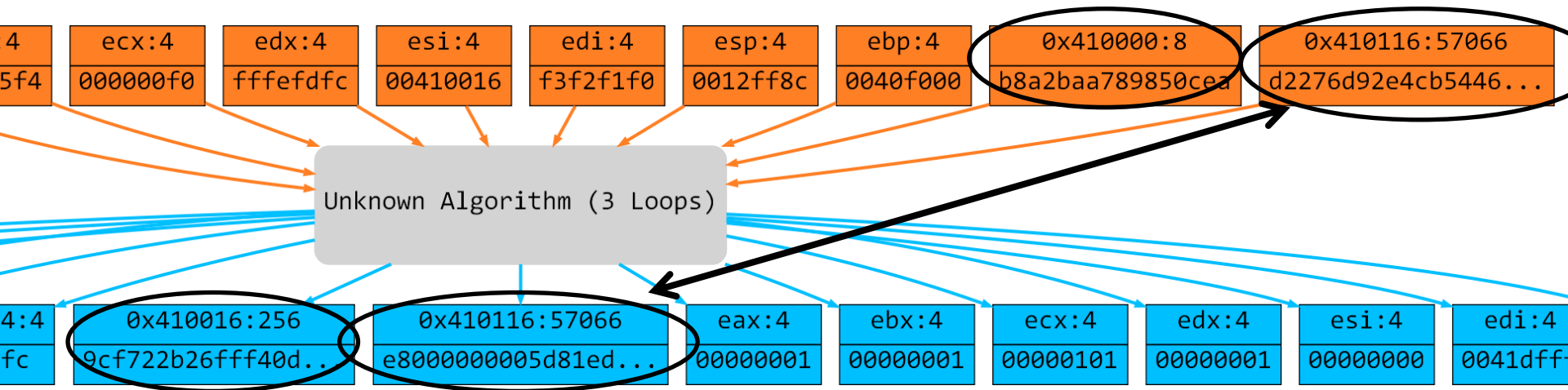
X86 ExecutableCode!

For the previous 3 loops, we extracted one algorithm:

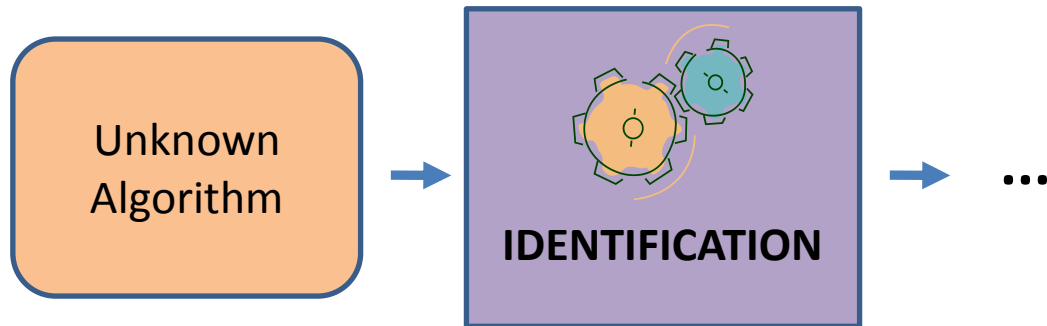


X86 ExecutableCode!

For the previous 3 loops, we extracted one algorithm:



X86 ExecutableCode!



```
STARTED AT
2012-04-10 16:29:05.135000
** Crypto Algorithm Identification starting !**
5 input parameters
4 output parameters
All possible input values generation... Done!
All possible output values generation... Done!
Build internal structure... Done!
Comparison phase starting... Test for RC4...
```

```
** Found RC4 **
```

```
==> Key (8 bytes) : b8a2baa789850cea
```

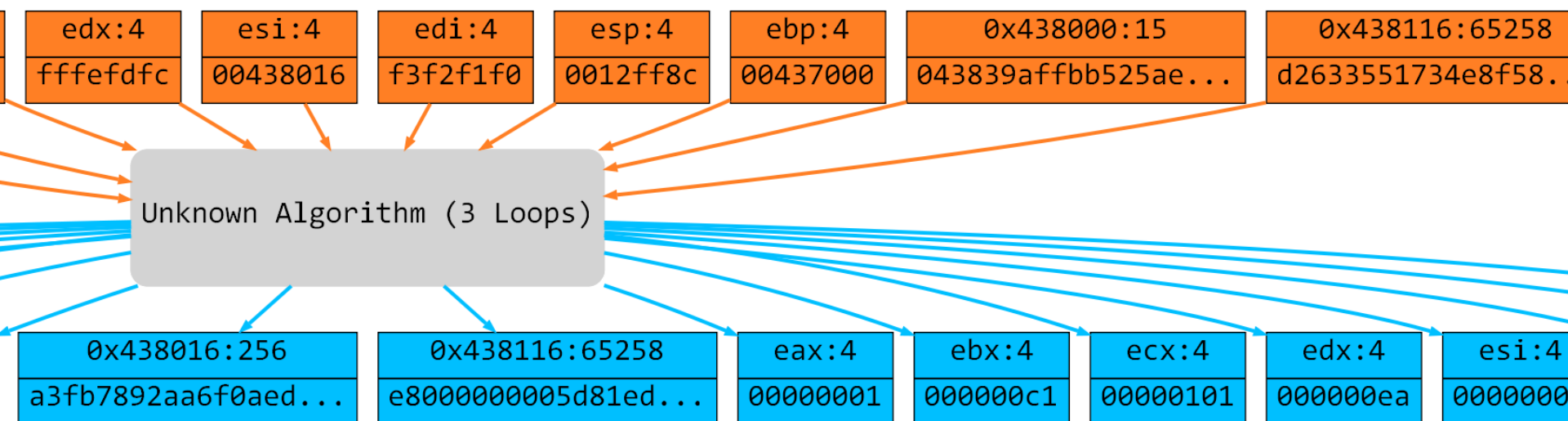
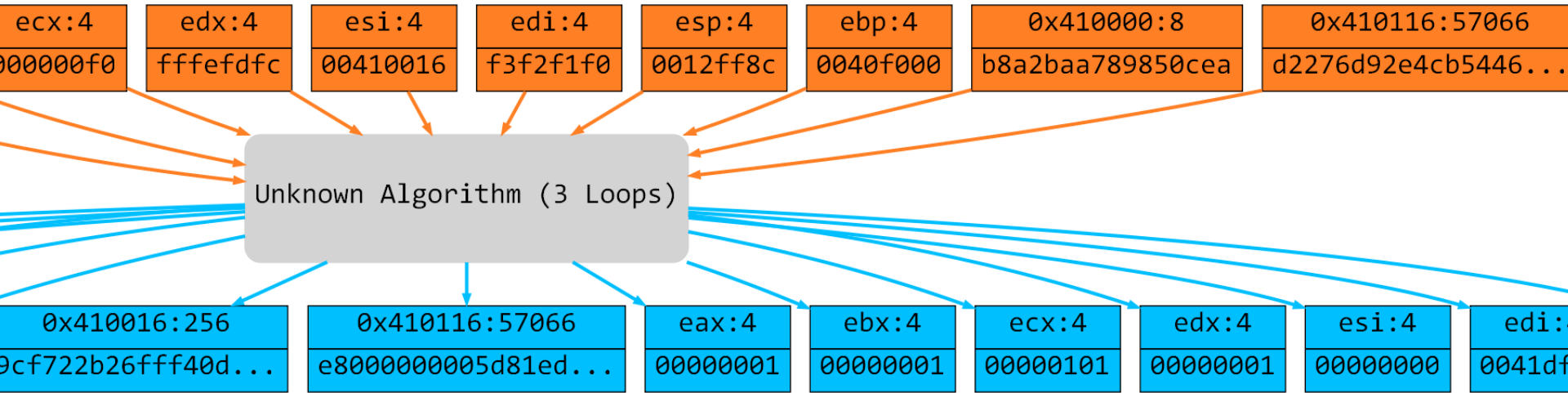
```
==> Crypted text (57066 bytes) : d2276d92e4cb5446...
```

```
==> Decrypted text (57066 bytes) : e8000000005d81ed...
```

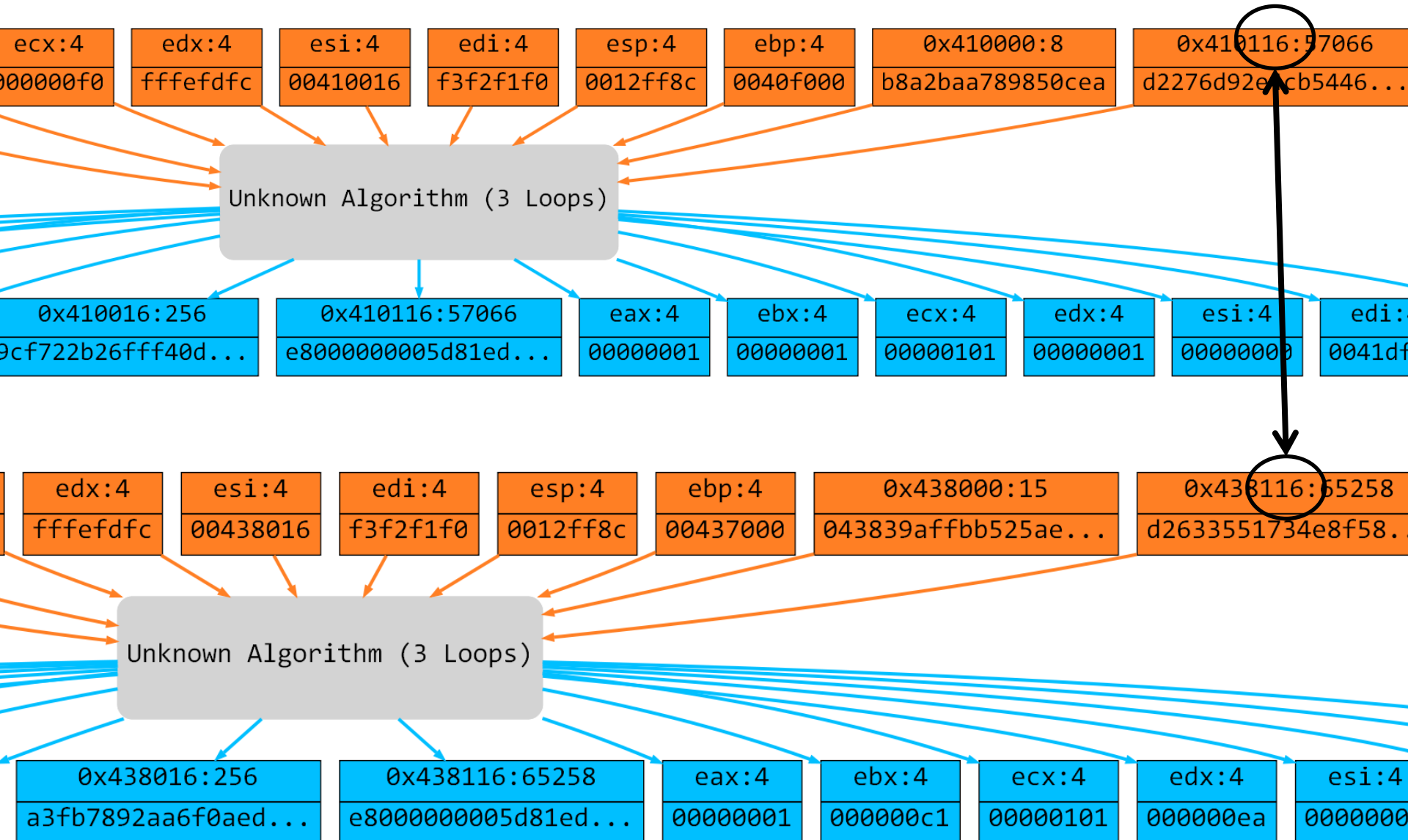
```
ENDED AT
2012-04-10 16:29:06.929000
```



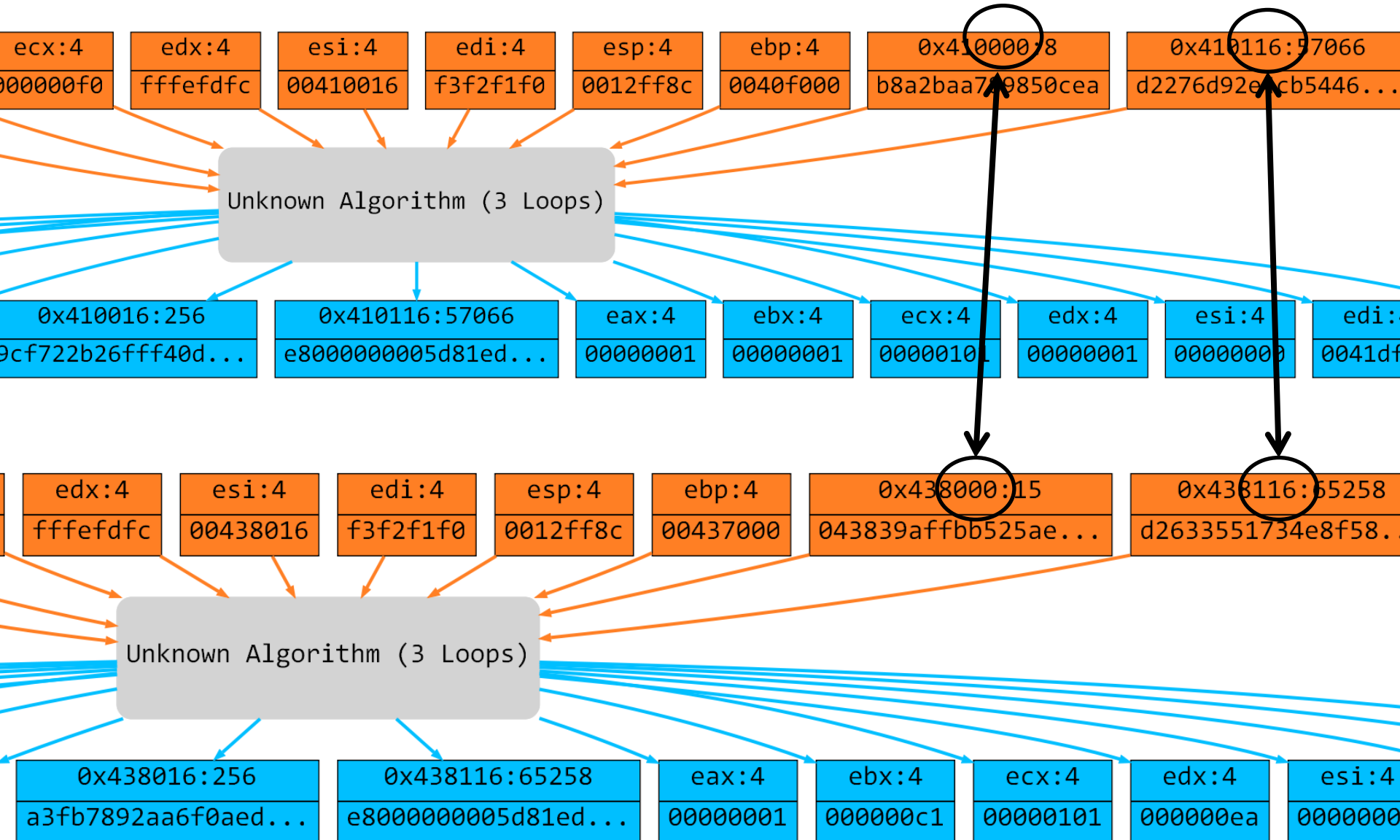
RC4 extracted from two Sality binaries



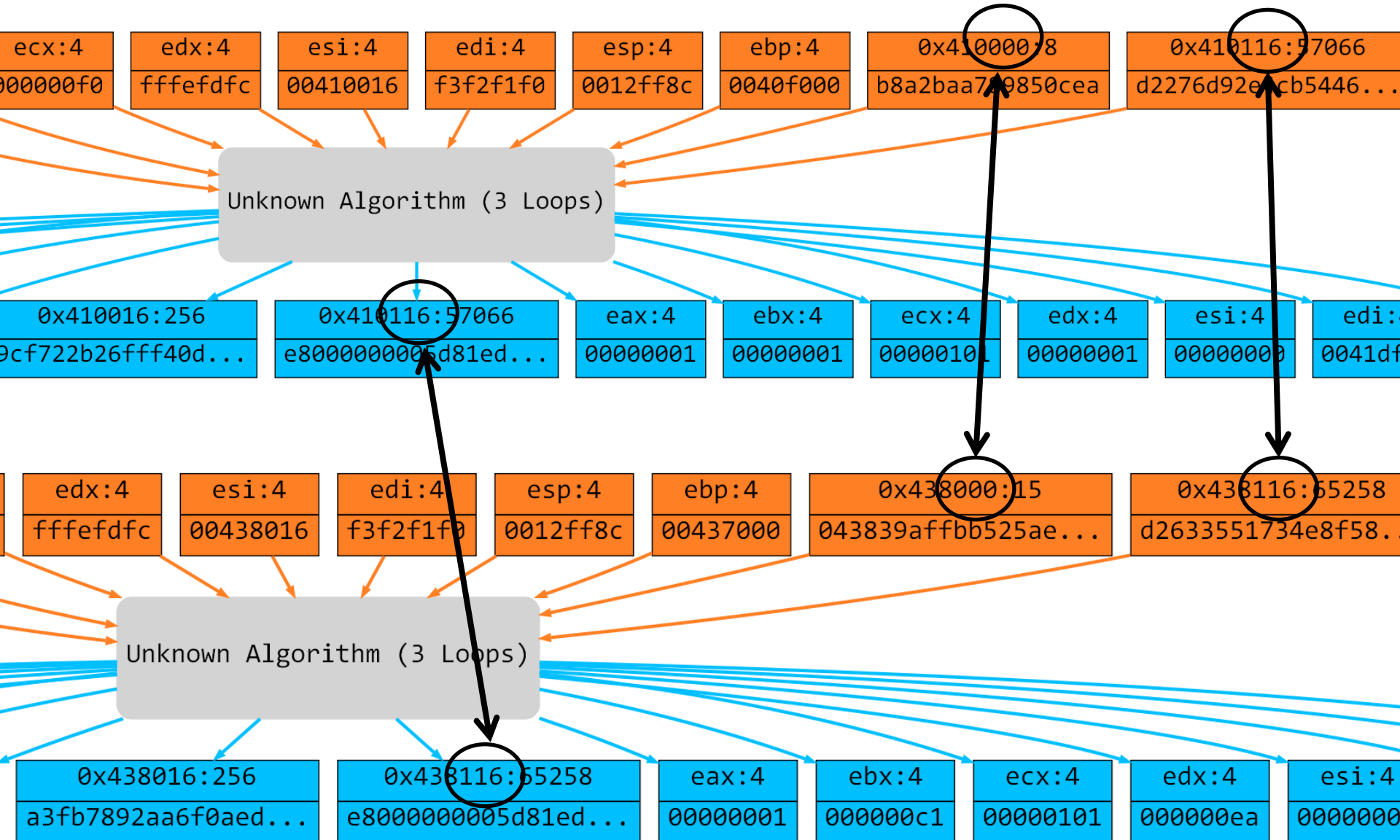
RC4 extracted from two Sality binaries



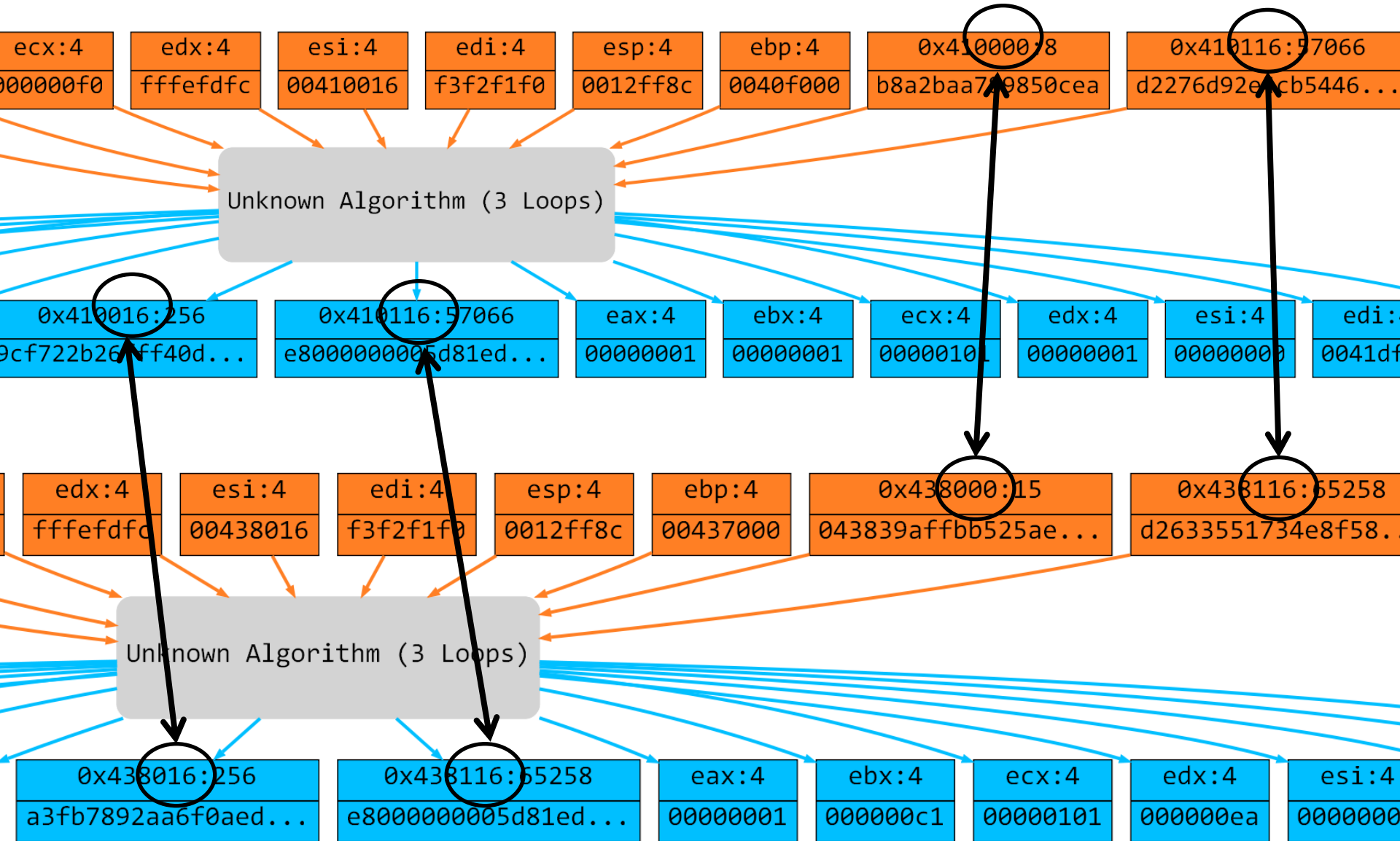
RC4 extracted from two Sality binaries



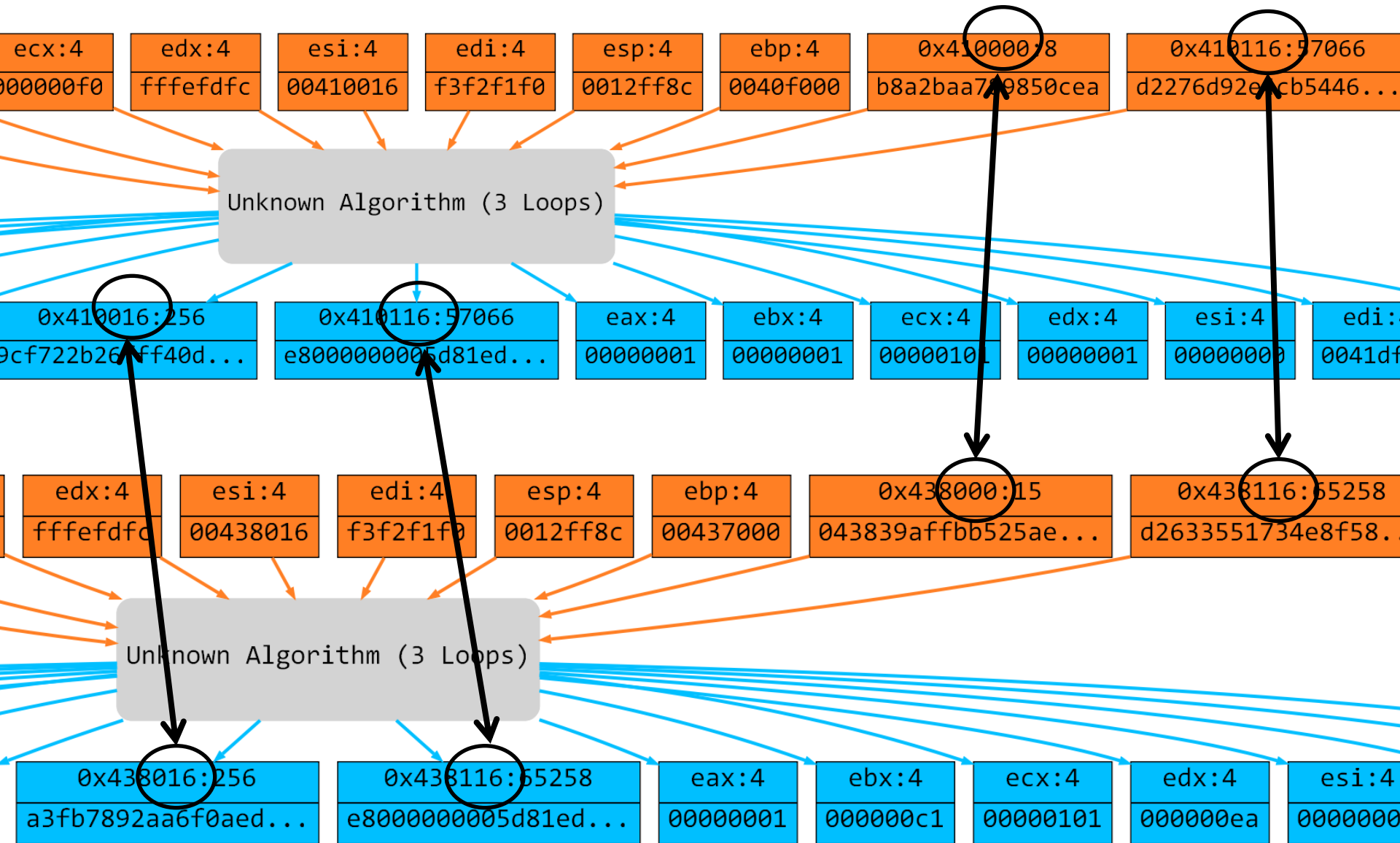
RC4 extracted from two Sality binaries



RC4 extracted from two Sality binaries



RC4 extracted from two Sality binaries



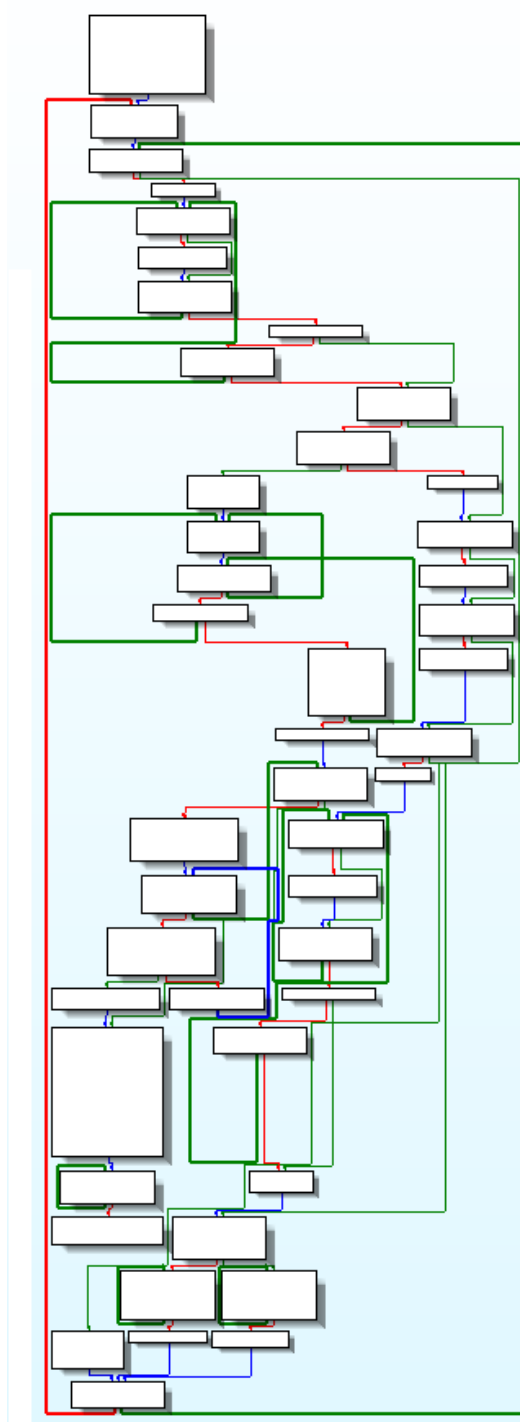
Crypto parameters always at the same offsets!

**EXAMPLE WHERE IT ~~DOESN'T~~
COULD WORK**

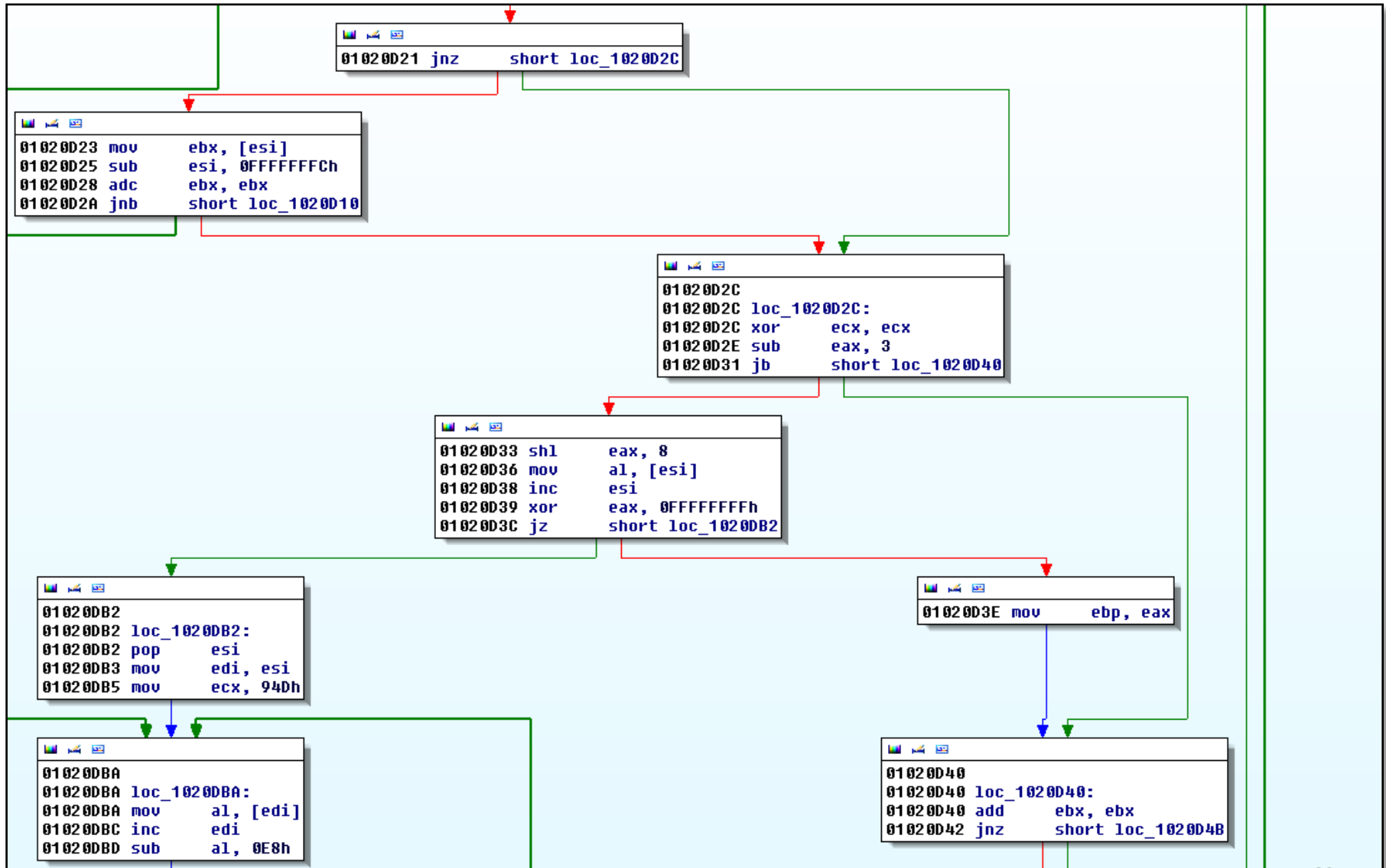
Compression Algorithms

- Compression algorithms have also a very particular input-output relationship.
- That should work too!

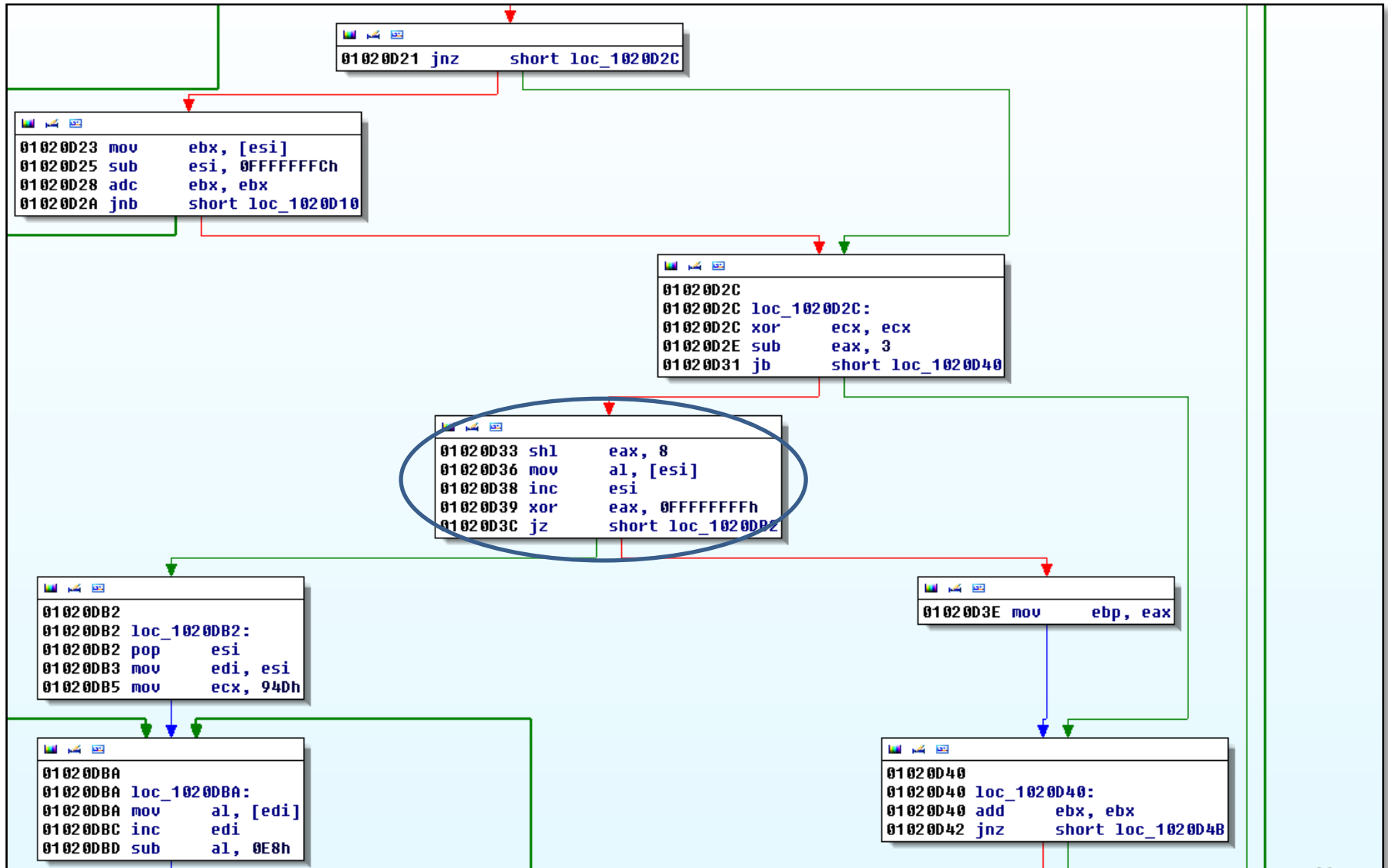
UPX LZMA



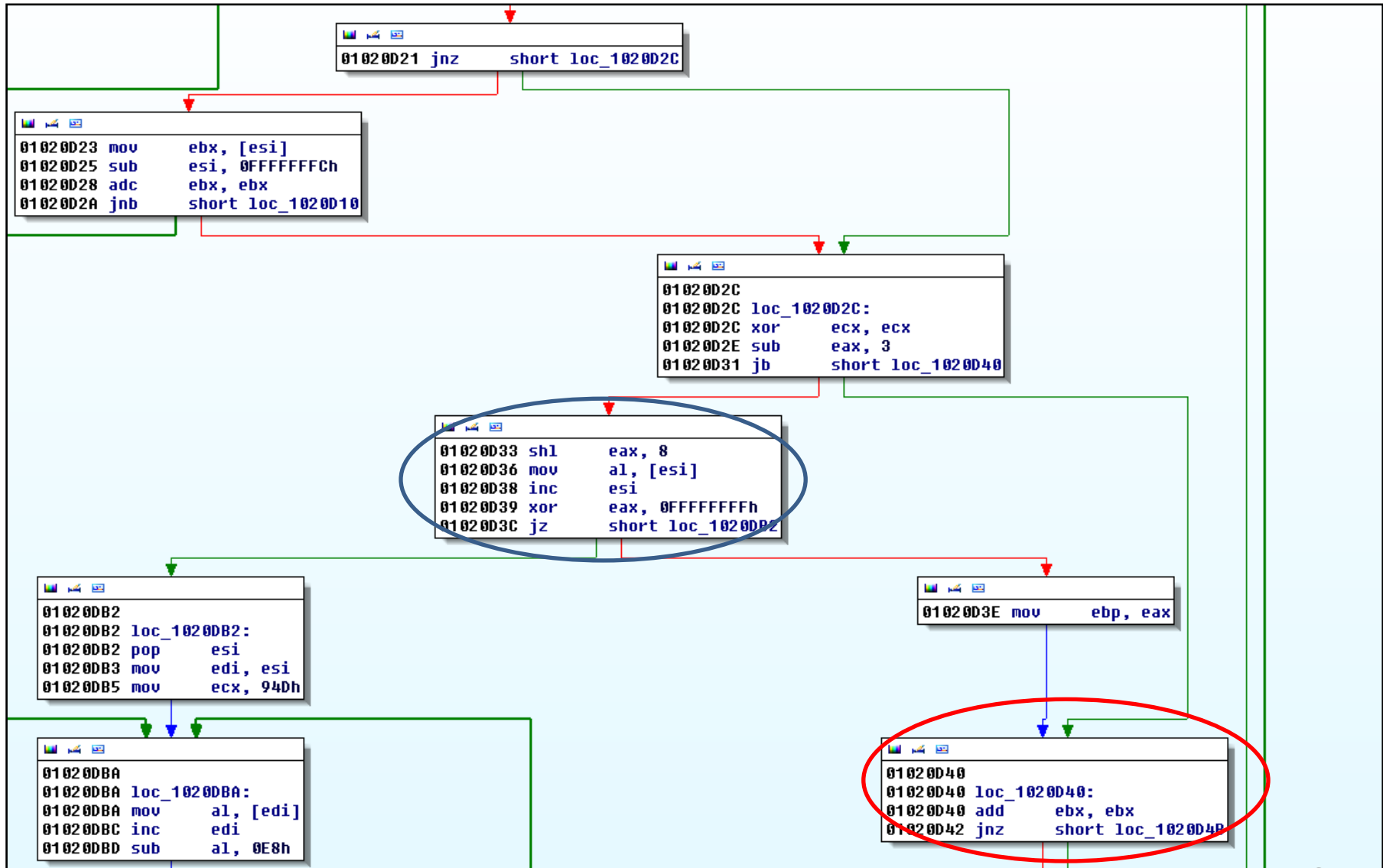
Fail... Why ?



Fail... Why ?



Fail... Why ?



Method Recap

1. We collect an execution trace.
2. We extract possible cryptographic algorithms with their parameter values.
3. We compare the input-output relationship with known algorithms.

We prove that the program behaves like a known crypto algorithm during one particular execution path.

Conclusion (1)

- Interesting alternative to syntactic-identification for crypto algorithms:
 - Resistance against usual obfuscation techniques.
 - Gives the exact parameters.
- Pure dynamic technique, you have to know how to exhibit interesting execution paths.
- It is easy to bypass, like any program analysis technique 😞

Conclusion (2)

- The identification process itself is generic:
 - Collect the execution trace
 - Extract the type of code you are looking for (here is the magic)
 - Get I/O values
 - Compare with reference implementations
- Nice work: Felix Gröbert **“Automatic Identification of Cryptographic Primitives in Software”**, 27th CCC

<http://code.google.com/p/kerckhoffs/>

What's Next ?

- That's only the beginning! Just wanted to show that it is feasible and useful.
- Extension to other crypto algorithms.
- Compression algorithms ? What should be the correct criteria to extract compression code ?
(my 2 cts: natural loops, that is back-edges on CFG)
- If we use several criteria to extract interesting code, how to combine them ?
- How to use the analyst knowledge ?
- Make a real tool. This one is just a PoC.

Thank you for your attention ;-)

Performances

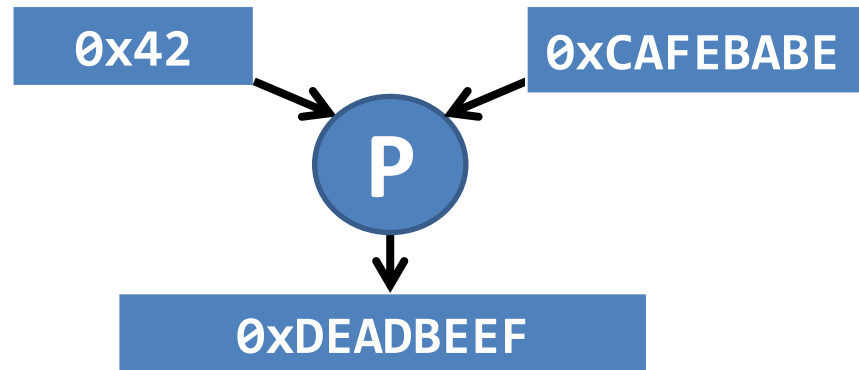
	Sality 1	Sality 2
Trace Size (instructions)	~1M	~4M
Time To Trace	5mn	10mn
Time To Extract Crypto Algoritm	4h	15h
Time To Identify	3mn	4mn

- The tool is just a PoC, no optimization **at all**.
- When the analysts knows where the algorithm is, it will reduce the trace size.

Answer: Brute-Force

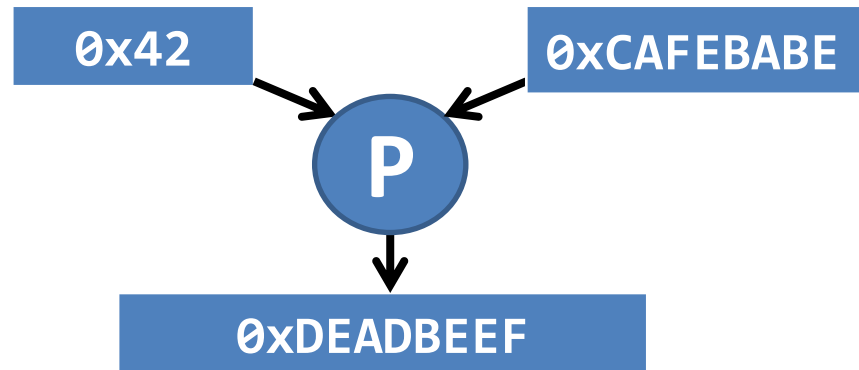
- We generate any possible input values by appending together **A** inputs, it gives a set **IN**.
- Same thing for output values, in the set **OUT**.
- For each reference implementation **F**:
 - We select the **IN** values that can fit into **F** parameters.
 - We execute **F** on each possible input value combination.
 - If **F** output is something in **OUT**, it is a success.

For example:



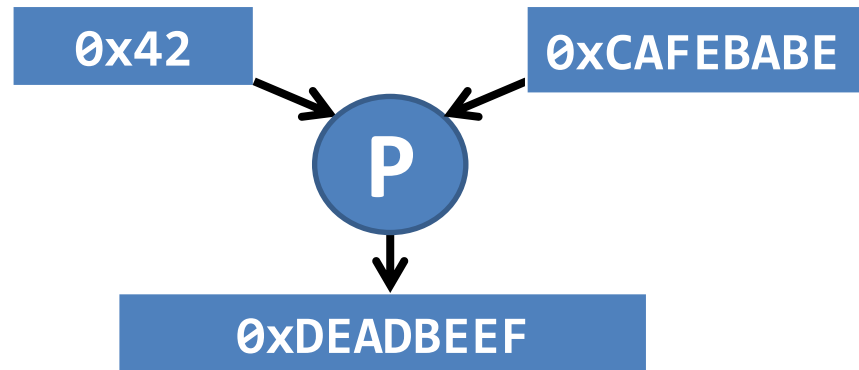
- For each known cryptographic algorithm **A**:

For example:



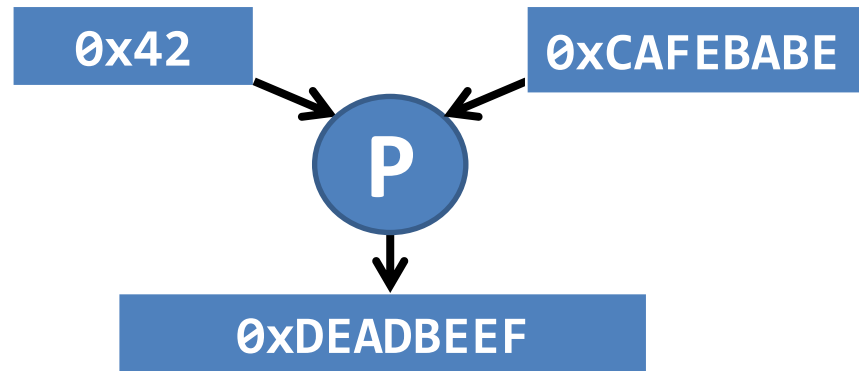
- For each known cryptographic algorithm **A**:
 - We execute a reference implementation of **A** on **0x42** and **0xCAFEBAE**.

For example:



- For each known cryptographic algorithm **A**:
 - We execute a reference implementation of **A** on **0x42** and **0xCAFEBAFE**.
 - We check if the output is **0xDEADBEEF**.

For example:



- For each known cryptographic algorithm **A**:
 - We execute a reference implementation of **A** on **0x42** and **0xCAFEBAE**.
 - We check if the output is **0xDEADBEEF**.
 - If so, **we have proved that P implements A on these particular input values.**